

THE RAMSEY NUMBER FOR HYPERGRAPH CYCLES I.

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ABSTRACT. Let \mathcal{C}_n denote the 3-uniform hypergraph *loose cycle*, that is the hypergraph with vertices v_1, \dots, v_n and edges $v_1v_2v_3, v_3v_4v_5, v_5v_6v_7, \dots, v_{n-1}v_nv_1$. We prove that every red-blue colouring of the edges of the complete 3-uniform hypergraph with N vertices contains a monochromatic copy of \mathcal{C}_n , where N is asymptotically equal to $5n/4$. Moreover this result is (asymptotically) best possible.

1. INTRODUCTION

A classical result in graph theory states that for $n \geq 5$, the 2-colour Ramsey number $r(C_n, C_n)$ of the cycle C_n with n vertices is $2n - 1$ if n is odd and $3n/2 - 1$ if n is even (Bondy and Erdős [1], Faudree and Schelp [3], Rosta [10]). Bondy and Erdős conjectured in 1973 that, for odd n , the 3-colour Ramsey number $r(C_n, C_n, C_n)$ equals $4n - 3$, and this appears to be quite a difficult problem. An asymptotic version was proved by Łuczak in 1999 [8], who showed that $r(C_n, C_n, C_n) = (4 + o(1))n$ if n is odd. Recently, Figaj and Łuczak [4] found the asymptotic value of the Ramsey number for all triples of even long cycles, proving, in particular, that $r(C_n, C_n, C_n) = (2 + o(1))n$ if n is even. Thus, in this case as well, the Ramsey number depends in a significant way on the parity of n .

In this paper we consider the problem of finding the 2-colour Ramsey number for a 3-uniform hypergraph cycle with n vertices. There are several natural definitions for a cycle in a 3-uniform hypergraph. The one we focus on here is what we call the *loose cycle* \mathcal{C}_n , which has vertices v_1, \dots, v_n and edge set $\{v_1v_2v_3, v_3v_4v_5, v_5v_6v_7, \dots, v_{n-1}v_nv_1\}$. Note that the number of vertices n of \mathcal{C}_n is necessarily even, and \mathcal{C}_n contains exactly $n/2$ vertices of degree one and $n/2$ vertices of degree two. The number of edges of \mathcal{C}_n , called the *length* of \mathcal{C}_n , is $n/2$. We remark that in another paper [7] we consider the analogous problem for the *tight cycle*, which has vertex set v_1, \dots, v_n and edge set $\{v_1v_2v_3, v_2v_3v_4, v_3v_4v_5, \dots, v_nv_1v_2\}$. The Ramsey number for the tight

cycle is larger than that of the loose cycle, and the tools needed for the argument are substantially different.

The Ramsey number $r(\mathcal{C}_n, \mathcal{C}_n)$ of \mathcal{C}_n is defined to be the smallest integer N for which every colouring of the edges of the complete 3-uniform hypergraph $K_N^{(3)}$ contains a monochromatic copy of \mathcal{C}_n , that is, a copy of \mathcal{C}_n whose edges are all coloured the same colour. We first note the following lower bound on $r(\mathcal{C}_n, \mathcal{C}_n)$.

Lemma 1.1. *We have $r(\mathcal{C}_{4k}, \mathcal{C}_{4k}) > 5k - 2$ and $r(\mathcal{C}_{4k+2}, \mathcal{C}_{4k+2}) > 5k + 1$.*

Proof. To prove the first assertion we exhibit a 2-colouring of the edges of $K_{5k-2}^{(3)}$ that contains no monochromatic copy of \mathcal{C}_{4k} . We partition the vertex set into parts A and B , where $|A| = k - 1$ and $|B| = 4k - 1$. We colour all edges that contain a vertex of A blue, and the rest red. Now this colouring cannot contain a red copy of \mathcal{C}_{4k} , since any such copy must have all vertices in B . Suppose that a blue copy of \mathcal{C}_{4k} exists, then each of its edges must contain a vertex of A , that is, a vertex-cover of \mathcal{C}_{4k} would be contained in A . But \mathcal{C}_{4k} has $2k$ edges and maximum degree two, and so cannot have a vertex cover of size smaller than k . Therefore, since $|A| = k - 1$, no blue copy of \mathcal{C}_{4k} can exist in this colouring, hence the lower bound is proved.

The argument for the second bound is exactly the same, except that we take $|A| = k$ and $|B| = 4k + 1$. \square

Our main aim in this paper is to prove that this lower bound is asymptotically tight, in other words that $r(\mathcal{C}_n, \mathcal{C}_n) \sim 5n/4$.

Theorem 1.2. *For all $\eta > 0$ there exists $n_0 = n_0(\eta)$ such that for every $n > n_0$, every 2-colouring of $K_{5(1+\eta)n/4}^{(3)}$ contains a monochromatic copy of \mathcal{C}_n .*

As in [8] and [4], our proof is based on the Regularity Lemma, here however we will need a hypergraph variant of this result. Note that in contrast to the graph case, the parity of the length of our cycle does not have a significant effect on its Ramsey number.

In our arguments we will often need to find loose *paths* between specified vertices. The *loose path* \mathcal{P}_n has vertices v_1, \dots, v_n and edge set $\{v_1v_2v_3, v_3v_4v_5, v_5v_6v_7, \dots, v_{n-2}v_{n-1}v_n\}$. By the *length* of a loose path we mean its number of edges, so \mathcal{P}_n has length $(n - 1)/2$.

2. PROOF OF THEOREM 1.2

In this section we state the principal lemmas required to prove Theorem 1.2, and show how they lead to the proof.

We begin by introducing a regularity lemma for hypergraphs. There are several generalizations of the Regularity Lemma of Szemerédi [11] to hypergraphs, due to various authors, e.g. [2], [5], [6], [9]. Here we will use the most straightforward one (see Chung [2]).

Let $\epsilon > 0$, let V_1, V_2, V_3 be disjoint vertex sets of size m , and let \mathcal{H} be a hypergraph such that each edge of \mathcal{H} contains exactly one element of each V_i for $i = 1, 2, 3$. Let $d = |\mathcal{H}|/m^3$. Then \mathcal{H} is said to be ϵ -regular of density d if for every choice of $X_i \subseteq V_i$ with $|X_1||X_2||X_3| > \epsilon m^3$ we have

$$\left| \frac{|\mathcal{H}[X_1, X_2, X_3]|}{|X_1||X_2||X_3|} - d \right| < \epsilon.$$

Here by $\mathcal{H}[X_1, X_2, X_3]$ we mean the subhypergraph of \mathcal{H} induced by the vertex set $X_1 \cup X_2 \cup X_3$. We often refer to the hypergraph $\mathcal{H} = \mathcal{H}[V_1, V_2, V_3]$ as an ϵ -regular *triple*. In this setting the (weak) regularity lemma for hypergraphs from [2] can be stated as follows.

Theorem 2.1. *Let $\epsilon > 0$ be given. Then there exists $t(\epsilon)$ such that every 3-uniform hypergraph \mathcal{M} has a partition $V_0 \cup V_1 \dots \cup V_t$ of its vertex set $V(\mathcal{M})$, where $t \leq t(\epsilon)$, such that*

- (1) $|V_0| < \epsilon N$, where $N = |V(\mathcal{M})|$,
- (2) $|V_i| = m$ for $i = 1, \dots, t$, where $(1 - \epsilon)N/t \leq m \leq N/t$,
- (3) all but at most ϵN^3 edges of \mathcal{M} lie in some ϵ -regular triple $\mathcal{M}[V_i, V_j, V_k]$ with $1 \leq i < j < k \leq t$.

We will prove Theorem 1.2 in three steps. In the first step, given a 2-colouring of the edges of $K_N^{(3)}$ for suitably chosen N , we apply Theorem 2.1 to the subhypergraph consisting of the red edges, to obtain a new 2-coloured structure called the *cluster* hypergraph. In the second step we show that the cluster hypergraph contains a monochromatic subhypergraph \mathcal{L} with certain special properties. Then in the third step, we prove that the monochromatic subhypergraph of $K_N^{(3)}$ corresponding to \mathcal{L} contains a copy of \mathcal{C}_n .

Step 1. Let η be given. We may assume $\eta < 1/5$. Let $\epsilon_0 < 1/5000$ be small enough such that $g(\epsilon_0) < 1 + \eta$, where $g(\epsilon_0) = (1 - 500\epsilon_0^{1/24})^{-3}$. We set $n_0 = 100t(\epsilon_0)^2\epsilon_0^{-1}$, where $t(\epsilon_0)$ is chosen in such a way that the assertion of Theorem 2.1 holds. For $n > n_0$, consider an arbitrary colouring of the edges of $K_N^{(3)}$ with red and blue, where $N = 5g(\epsilon_0)n/4$. Let \mathcal{J} denote the subhypergraph consisting of the edges coloured red, and let $\bar{\mathcal{J}}$ denote the blue subhypergraph. We apply Theorem 2.1 to \mathcal{J} with ϵ_0 to obtain a partition $V_0 \cup V_1 \dots \cup V_t$ of $V(K_N^{(3)})$ with the properties stated in Theorem 2.1. Then we have the following lemma, whose very standard proof appears in Section 3.

Lemma 2.2. *Let $\epsilon = \epsilon_0^{1/4}$. Then*

- (1) *all but at most ϵt^3 triples $\mathcal{J}[V_i, V_j, V_k]$ are ϵ -regular,*
- (2) *if $\mathcal{J}[V_i, V_j, V_k]$ is ϵ -regular with density d then $\bar{\mathcal{J}}[V_i, V_j, V_k]$ is ϵ -regular with density $1 - d$,*
- (3) *$20t < \epsilon m$.*

We define the *cluster hypergraph* \mathcal{J}_0 as follows. The vertex set of \mathcal{J}_0 is $\{1, \dots, t\}$, and ijk forms an edge of \mathcal{J}_0 precisely when $\mathcal{J}[V_i, V_j, V_k]$ is ϵ -regular (of some density). Then by Lemma 2.2, all but at most ϵt^3 of the triples ijk are edges of \mathcal{J}_0 .

We consider the edges of \mathcal{J}_0 to be coloured red or blue, as follows. We colour the edge $ijk \in \mathcal{J}_0$ red if $\mathcal{J}[V_i, V_j, V_k]$ has density at least $1/2$, otherwise we colour it blue. Then by Lemma 2.2(2), if ijk is coloured blue, then $\bar{\mathcal{J}}[V_i, V_j, V_k]$ is ϵ -regular of density more than $1/2$. This completes Step 1.

Step 2. To describe Step 2 we need to introduce a few more definitions. Let \mathcal{M} be a 3-uniform hypergraph. The *shadow graph* $\Gamma(\mathcal{M})$ of \mathcal{M} is defined on the vertex set $V(\mathcal{M})$ by joining vertices x and y by an edge if and only if there exists an edge $xyz \in \mathcal{M}$. We call \mathcal{M} a *connected* hypergraph if $\Gamma(\mathcal{M})$ is connected in the ordinary graph sense. A subhypergraph of \mathcal{M} that is maximal with respect to being connected is called a *component* of \mathcal{M} . Thus a component of \mathcal{M} is determined by its set of vertices. We will use this concept when \mathcal{M} is the subhypergraph of a 2-coloured hypergraph \mathcal{K} consisting of all the red edges (or all the blue edges). A component of \mathcal{M} will be called a *monochromatic component* of \mathcal{K} , so each monochromatic component is either red or blue. Thus each vertex in a 2-coloured hypergraph \mathcal{K} is in one red monochromatic component and one blue monochromatic component.

We also define a special small hypergraph called a *diamond*. A diamond \mathcal{D} has vertex set $\{x_1, x_2, x_3, x_4\}$ and edge set $\{x_1x_2x_3, x_2x_3x_4\}$. The two vertices x_2 and x_3 of degree 2 in \mathcal{D} are called the *central points* of \mathcal{D} .

Step 2 of the proof of Theorem 1.2 is accomplished by the following lemma.

Lemma 2.3. *The cluster hypergraph \mathcal{J}_0 has a monochromatic component \mathcal{L} that contains $s = \lceil f(\epsilon) \frac{n}{N} \frac{t}{4} \rceil$ vertex-disjoint diamonds $\mathcal{D}_1, \dots, \mathcal{D}_s$, where $f(\epsilon) = (1 - 500\epsilon^{1/6})^{-2}$.*

The proof of Lemma 2.3 appears in Section 5. Note that since $N = 5g(\epsilon_0)n/4$, and $f(\epsilon)$ and $g(\epsilon_0)$ are very close to 1, the parameter s is approximately equal to $t/5$.

Step 3. To complete Step 3 we show that the monochromatic component \mathcal{L} guaranteed by Lemma 2.3 corresponds to a monochromatic subhypergraph \mathcal{H} of $\mathcal{J}^* = \mathcal{J}$ or $\bar{\mathcal{J}}$ that contains a copy of \mathcal{C}_n . If \mathcal{L} is red then we let $\mathcal{J}^* = \mathcal{J}$, if \mathcal{L} is blue then we choose $\mathcal{J}^* = \bar{\mathcal{J}}$. Then \mathcal{H} is defined in the natural way, that is, $V(\mathcal{H}) = \bigcup\{V_i : i \in V(\mathcal{L})\}$ and $\mathcal{H} = \bigcup\{\mathcal{J}^*[V_i, V_j, V_k] : ijk \in \mathcal{L}\}$.

Lemma 2.4. \mathcal{H} contains a copy of \mathcal{C}_n .

The proof of Lemma 2.4 appears in Section 4.

We remark in passing that it may seem more natural to look for a monochromatic component in the cluster hypergraph that contains a large matching of edges, rather than a large set of disjoint diamonds. However, we are not guaranteed that a large enough component of this type exists in every colouring. We would need the matching to be of size at least $4t/15$ in order to find a copy of \mathcal{C}_n in the corresponding subhypergraph of \mathcal{H} . However, there exist colourings of the cluster hypergraph for which no such matching exists in a monochromatic component. For example, if we colour all edges incident to a set A of size $t/4$ blue, and the rest red, the largest monochromatic matching has size only $t/4$.

3. PROOF OF LEMMA 2.2

With the definitions given in Step 1, we consider the first assertion of Lemma 2.2. We begin by claiming that if $|\mathcal{J}[V_i, V_j, V_k]| < \epsilon_0^{1/2} m^3$, then it is ϵ -regular. To see this, note that $0 \leq d < \epsilon_0^{1/2} < \epsilon$ for this triple, so if $|X_i||X_j||X_k| > \epsilon m^3$, then

$$0 \leq \frac{|\mathcal{J}[X_i, X_j, X_k]|}{|X_i||X_j||X_k|} < \frac{\epsilon_0^{1/2} m^3}{\epsilon_0^{1/4} m^3} = \epsilon.$$

Therefore the density of $\mathcal{J}[X_i, X_j, X_k]$ differs from d by less than ϵ , as required.

Suppose on the contrary that at least ϵt^3 triples are not ϵ -regular in \mathcal{J} . Then they are certainly not ϵ_0 -regular either. By the above claim, the total number of edges of \mathcal{J} contained in these triples is at least $\epsilon \epsilon_0^{1/2} m^3 t^3 > \epsilon_0^{3/4} (1 - \epsilon_0)^3 N^3$ by Theorem 2.1(2). But since $\epsilon_0 < 1/5000$ this number is more than $\epsilon_0 N^3$, contradicting Theorem 2.1(3). Therefore the first assertion of Lemma 2.2 holds.

To check the second assertion, note that $|\bar{\mathcal{J}}[X_i, X_j, X_k]| = |X_i||X_j||X_k| - |\mathcal{J}[X_i, X_j, X_k]|$. In particular, the densities $d_{\bar{\mathcal{J}}}$ of $|\bar{\mathcal{J}}[X_i, X_j, X_k]|$ and $d_{\mathcal{J}}$ of $|\mathcal{J}[X_i, X_j, X_k]|$ are related by $d_{\bar{\mathcal{J}}} = 1 - d_{\mathcal{J}}$. Thus if $\mathcal{J}[V_i, V_j, V_k]$

is ϵ -regular then for any sets $X_i \subseteq V_i$, $X_j \subseteq V_j$ and $X_k \subseteq V_k$ with $|X_i||X_j||X_k| > \epsilon m^3$ we have

$$\frac{|\bar{\mathcal{J}}[X_i, X_j, X_k]|}{|X_i||X_j||X_k|} = 1 - \frac{|\mathcal{J}[X_i, X_j, X_k]|}{|X_i||X_j||X_k|},$$

and so this number differs from $1 - d_{\mathcal{J}}$ by less than ϵ as required.

The third assertion is simply a calculation from the definitions given in Step 1. We have $N > n > \frac{100t(\epsilon_0)^2}{\epsilon_0} > \frac{100t^2}{\epsilon}$ so $\epsilon(1 - \epsilon)N > 20t^2$. This together with Theorem 2.1(2) implies $\epsilon m > \epsilon(1 - \epsilon_0)\frac{N}{t} > \epsilon(1 - \epsilon)\frac{N}{t} > 20t$ as required.

4. PROOF OF LEMMA 2.4

To prove Lemma 2.4, we will use the following strategy to find the long cycle \mathcal{C}_n in the hypergraph \mathcal{H} defined in Step 3. We first trace a “route” in the monochromatic component \mathcal{L} of the cluster hypergraph \mathcal{J}_0 , that visits all of the s disjoint diamonds $\mathcal{D}_1, \dots, \mathcal{D}_s$ in \mathcal{L} (Lemma 4.1). Then, using Lemma 4.3, we choose a collection of short loose paths (of length three or six) in the hypergraph \mathcal{H} , that link together to form a cycle, following the chosen route. Finally, to obtain the cycle \mathcal{C}_n in \mathcal{H} , for each diamond $\mathcal{D}_i = \{hjk, hjp\}$, we replace the short path that starts in V_h and ends in V_j by a long path with the same end-vertices, that uses almost all the vertices in $V_h \cup V_j \cup V_k \cup V_p$. This step uses Lemma 4.5.

Let $G = \Gamma(\mathcal{L})$ be the shadow graph of \mathcal{L} , as defined in Step 2. Then for each of the s disjoint diamonds $\mathcal{D}_i = \{hjk, hjp\}$ in \mathcal{L} , there is an edge hj of G joining the two central points of \mathcal{D}_i . Let M denote the matching in G consisting of these s disjoint edges. By a *closed directed trail* in G we mean a sequence of (not necessarily distinct) vertices $x_1, x_2, \dots, x_r = x_1$ of G , together with the edges (x_i, x_{i+1}) for $i = 1, \dots, r - 1$ which we consider as being directed from x_i to x_{i+1} and call *arcs*.

The following easy lemma provides the route that our cycle will eventually follow.

Lemma 4.1. *Let $G = \Gamma(\mathcal{L})$, and let M be defined as above. Then G contains a closed directed trail with $r \leq 2t$ vertices that uses every edge of M as an arc, and no edge of G is used as an arc more than once in each direction.*

Proof. Recall that $|V(G)| \leq t$. Let T be any spanning tree of G that contains the edges of M . We can easily obtain a closed directed trail with the desired properties as follows. Fix a planar embedding of T , and

let x_1 be the root of T . Construct the trail by taking a walk starting at x_1 around the boundary of the (single) face of the embedding, ending back at x_1 . Then each edge of T is traversed exactly once in each direction, so the number of vertices of this trail is at most $2t$. Moreover, every edge of M is used as an arc. \square

For each diamond \mathcal{D}_i , $1 \leq i \leq s$, we choose one arc that uses the edge of M corresponding to \mathcal{D}_i and call it the *diamond arc* for \mathcal{D}_i .

When choosing short paths that will link up to form a cycle, we need to make sure the end-vertices of these paths can indeed coincide. Let V_h, V_j , and V_k be any vertex classes such that $\mathcal{H}[V_h, V_j, V_k]$ is ϵ -regular with density $d > 2\epsilon^{1/2}$ (recall $|V_h| = |V_j| = |V_k| = m$), and for $i \in \{h, j, k\}$ let $U_i \subseteq V_i$ be subsets. We say that a vertex $x \in V_h$ is *good* for the triple $\mathcal{H}[U_h, U_j, U_k]$ if

- (i) for at least $d|U_j|/2$ vertices $y \in U_j$, there are at least $d|U_k|/2$ vertices $z \in U_k$ such that $xyz \in \mathcal{H}$, and
- (ii) for at least $d|U_k|/2$ vertices $z \in U_k$, there are at least $d|U_j|/2$ vertices $y \in U_j$ such that $xyz \in \mathcal{H}$.

We define vertices in V_j and V_k to be good for $\mathcal{H}[U_h, U_j, U_k]$ in a similar way. The next lemma implies that, as long as U_h, U_j and U_k are reasonably large, most vertices in $V_h \cup V_j \cup V_k$ are good.

Lemma 4.2. *Suppose $d > 2\epsilon^{1/2}$. With the above definitions, the number of vertices in V_h that are good for $\mathcal{H}[U_h, U_j, U_k]$ is at least $m - \frac{\epsilon m^3}{|U_j||U_k|}$. The analogous bounds hold for V_j and V_k . In particular, for each $i \in \{h, j, k\}$ the number of vertices in V_i that are good for $\mathcal{H}[V_h, V_j, V_k]$ is at least $(1 - \epsilon)m$.*

Proof. Suppose $x \in V_h$ fails to satisfy (i). Then the total number of edges of $\mathcal{H}[\{x\}, U_j, U_k]$ is smaller than

$$(d|U_j|/2)|U_k| + (1 - d/2)|U_j|(d|U_k|/2) = (d - d^2/4)|U_j||U_k|.$$

If x fails to satisfy (ii) then we find the same bound. Let $X \subset V_h$ denote the set of vertices that are not good. Then $|\mathcal{H}[X, U_j, U_k]| \leq |X|(d - d^2/4)|U_j||U_k|$. Thus if $|X||U_j||U_k| > \epsilon m^3$ we find

$$\frac{|\mathcal{H}[X, U_j, U_k]|}{|X||U_j||U_k|} \leq d - d^2/4 < d - \epsilon,$$

which contradicts the fact that $\mathcal{H}[V_h, V_j, V_k]$ is ϵ -regular. Therefore the number of vertices in V_h that are not good for $\mathcal{H}[U_h, U_j, U_k]$ is at most $\frac{\epsilon m^3}{|U_j||U_k|}$. Similar bounds can be found for V_j and V_k . \square

The set of vertices in $V_h \cup V_j \cup V_k$ that are good for $\mathcal{H}[V_h, V_j, V_k]$ will simply be called *good*. To make sure the end-vertices of our short paths coincide where necessary, we will require them to be good. This idea will be made precise in the upcoming proof of Lemma 2.4.

We will also need to choose our short paths in ϵ -regular triples so that they avoid some “bad” set B of vertices. The next lemma allows us to find these short paths (of length three).

Lemma 4.3. *Let $\mathcal{H}[V_h, V_j, V_k]$ be an ϵ -regular triple with density $d > 2\epsilon^{1/3}$. Then for every pair of good vertices $x \in V_h$ and $y \in V_j$, and for every set $B \subset V_h \cup V_j \cup V_k \setminus \{x, y\}$ that contains all non-good vertices and satisfies $|B \cap V_i| < (d/2 - \epsilon^{1/3})m$ for $i \in \{h, j, k\}$, there is a path of length three in $\mathcal{H}[V_h, V_j, V_k]$ joining x to y that is disjoint from B (and hence contains only good vertices). Moreover the path can be chosen so that one vertex of degree two in the path is in V_h , and the other is in V_j .*

Proof. Since x is good, there exists a set $U_x \subset V_j$, $|U_x| \geq dm/2$, such that for each $z \in U_x$, there are at least $dm/2$ vertices $w \in V_k$ such that $xzw \in \mathcal{H}$. Since y is good, there exists a set $U_y \subset V_h$ with similar properties. Then, writing $b = (d/2 - \epsilon^{1/3})m$, we have

$$|U_x \setminus B||U_y \setminus B||V_k \setminus B| > (d/2 - b)^2(1 - b)m^3 > \epsilon m^3.$$

Therefore, since \mathcal{H} is ϵ -regular, we know that $|\mathcal{H}[U_x \setminus B, U_y \setminus B, V_k \setminus B]| \geq (d - \epsilon)|U_x \setminus B||U_y \setminus B||V_k \setminus B|$. We may therefore choose vertices $z \in U_x \setminus \{y\}$ and $v \in U_y \setminus \{x\}$, and distinct good vertices $u_h, u_j, u_k \in V_k$ such that $xzu_h, yvu_j, zvu_k \in \mathcal{H}$. This gives the required path. \square

The following immediate consequence of the above lemma will be used to adjust the parity of the length of a path.

Corollary 4.4. *Let $\mathcal{H} = \mathcal{H}[V_h, V_j, V_k]$ be an ϵ -regular triple with density $d > 2\epsilon^{1/3}$. Then for every pair of good vertices $x \in V_h$ and $y \in V_j$, and for every set $B \subset V_h \cup V_j \cup V_k \setminus \{x, y\}$ that contains all non-good vertices and satisfies $|B \cap V_i| < (d/2 - \epsilon^{1/3})m - 6$ for each $i \in \{h, j, k\}$, there is a path of length six in \mathcal{H} joining x to y that is disjoint from B (and hence contains only good vertices).*

Proof. Let w be any vertex in $V_k \setminus B$. Then by Lemma 4.3 there is a path \mathcal{P} of length three joining x and w that avoids $B \cup \{y\}$. Let $B' = B \cup V(\mathcal{P}) \setminus \{w\}$. Then applying Lemma 4.3 again with w, y and B' we may extend \mathcal{P} to a path of length six with the required properties. \square

Our final preparatory lemma tells us that in a “blown-up diamond” we can find long paths between specified vertices, that avoid any given small set B . We remark that the condition on the parity of the path is not essential, it is present simply to make the proof a bit easier.

Lemma 4.5. *Let V_h, V_j, V_k, V_p be classes such that $\mathcal{H}_1 = \mathcal{H}[V_h, V_j, V_k]$ and $\mathcal{H}_2 = \mathcal{H}[V_j, V_k, V_p]$ are both ϵ -regular of density $d \geq 32\epsilon^{1/6}$. Let $x \in V_j$ and $y \in V_k$ be good vertices for both \mathcal{H}_1 and \mathcal{H}_2 . Let $B \subset V_h \cup V_j \cup V_k \cup V_p \setminus \{x, y\}$ be such that B contains all vertices that are not good for \mathcal{H}_1 or not good for \mathcal{H}_2 , and $|B \cap V_i| < 2\epsilon m$ for each $i \in \{h, j, k, p\}$. Then there is a path joining x and y in $\mathcal{H}_1 \cup \mathcal{H}_2$ of length ℓ for all odd integers ℓ satisfying $2dm/7 \leq \ell \leq (1 - 2\epsilon^{1/6})2m$ that avoids B .*

Proof. We prove the lemma by means of the following two claims. The first is an easy greedy argument using Lemma 4.3.

Claim 4.6. *Let $0 \leq q < \lfloor (d/2 - \epsilon^{1/3})m/6 \rfloor$, and let $B' \subset V_h \cup V_j \cup V_k \setminus \{x, y\}$ be such that B' contains all vertices that are not good for \mathcal{H}_1 , and $|B' \cap V_i| < (d/2 - \epsilon^{1/3})m - 6q$ for each $i \in \{h, j, k\}$. Then there is a path joining x and y in \mathcal{H}_1 of length $6q + 3$ that avoids B' , whose degree-two vertices all lie in V_j or V_k .*

Proof. We prove the claim by induction on q . For $q = 0$ it is immediately given by Lemma 4.3. Assume $q \geq 1$ and that the claim is true for smaller values of q . Let B' be given. Choose $z \in V_k \setminus (B' \cup \{y\})$, then z is good. Then by the induction hypothesis, there is a path \mathcal{P} of length $6(q-1) + 3$ joining x to z that avoids $B' \cup \{y\}$, and whose degree-two vertices all lie in V_j or V_k . Note then $|V(\mathcal{P}) \cap V_h| = 6(q-1) + 3$ and $|V(\mathcal{P}) \cap V_j| = |V(\mathcal{P}) \cap V_k| = 3(q-1) + 2$.

Let $B'' = B' \cup V(\mathcal{P}) \cup \{y\} \setminus \{z\}$. Then $|B'' \cap V_i| < (d/2 - \epsilon^{1/3})m - 3$. Let $w \in V_j \setminus B''$. Then w is good, so by Lemma 4.3 there is a path \mathcal{P}' of length three joining z and w that avoids B'' . Finally let $B''' = B'' \cup V(\mathcal{P}') \setminus \{w, y\}$. Then $|B''' \cap V_i| < (d/2 - \epsilon^{1/3})m$. Apply Lemma 4.3 once more to obtain a path \mathcal{P}'' of length 3 joining w to y that avoids B''' . The concatenation $\mathcal{P}\mathcal{P}'\mathcal{P}''$ then has length $6q + 3$, avoids B' , and all its degree-two vertices lie in $V_j \cup V_k$. \square

Our second claim states that any reasonably long path can be lengthened by two, unless it uses almost all the vertices of $\mathcal{H}_1 \cup \mathcal{H}_2$.

Claim 4.7. *With the assumptions of Lemma 4.5, let \mathcal{P} be a path in $\mathcal{H}_1 \cup \mathcal{H}_2$ joining x and y of length ℓ , where $dm/4 < \ell < (1 - 2\epsilon^{1/6})2m$, that avoids B . Suppose all degree-two vertices of \mathcal{P} lie in $V_j \cup V_k$.*

Then there exists a path \mathcal{P}' joining x and y that avoids B , all of whose degree-two vertices lie in $V_j \cup V_k$, of length $\ell + 2$.

Proof. The aim is to replace one edge of \mathcal{P} by three edges, however we must choose the edge to be replaced appropriately.

For each $i \in \{h, j, k, p\}$ let $U_i \subset V_i$ denote the available vertices in V_i , that is $U_i = V_i \setminus (V(\mathcal{P}) \cup B)$. Then by the condition on the length of \mathcal{P} , we have $|U_i| > \epsilon^{1/6}m$ for $i \in \{j, k\}$, and $|U_i| > \epsilon^{1/6}m$ for some $i \in \{h, p\}$, say without loss of generality h .

Now by Lemma 4.2, for each $i \in \{h, j, k\}$ at most $\epsilon^{2/3}m$ vertices of $V_i \cap V(\mathcal{P})$ are not good for $\mathcal{H}[U_h, U_j, U_k]$. Therefore, since $|V_i \cap V(\mathcal{P})| \geq \ell/2 \geq dm/8 > 4\epsilon^{2/3}m$, there exist $x_1 \in V_j \cap V(\mathcal{P})$ and $y_1 \in V_k \cap V(\mathcal{P})$ that are both good for $\mathcal{H}[U_h, U_j, U_k]$, such that $x_1 y_1 z_1$ is an edge of \mathcal{P} for some $z_1 \in V_h$. Let $U_k(x_1) \subset U_k$ be the set guaranteed by the goodness of x_1 , so $|U_k(x_1)| \geq d|U_k|/2 > \epsilon^{1/3}m$ has property (i). Let $U_j(y_1) \subset U_j$ be defined similarly (having property (ii)). Then we have $|U_h||U_j(y_1)||U_k(x_1)| > \epsilon m^3$, and so there exist $u \in U_k(x_1)$ and $w \in U_j(y_1)$ such that $uwz \in \mathcal{H}$ for some $z \in U_h$. Thus, a path of length three from x_1 to y_1 can easily be found, that is disjoint from the rest of \mathcal{P} and from B , and whose degree-two vertices lie in V_j and V_k . \square

The proof of Lemma 4.5 now follows immediately, as we can find a path of length $\ell_0 = 6(\lfloor (d/2 - \epsilon^{1/3})m/6 \rfloor - 1) + 3 > dm/4$ from x to y using Claim 4.6, and then extend it step by step using Claim 4.7 until a path of the desired length is reached. \square

We are now ready to prove Lemma 2.4.

Proof of Lemma 2.4. Let \mathcal{L} be as in Lemma 2.3, and let $T = x_1 x_2 \dots x_r = x_1$ be the closed directed trail in $\Gamma(\mathcal{L})$ given by Lemma 4.1. For each arc (x_{i-1}, x_i) , choose an ϵ -regular triple $\mathcal{H}_i = \mathcal{H}[V_{x_{i-1}}, V_{x_i}, V_{k_i}]$ of density at least $1/2$ for some k_i (the existence of which is guaranteed by the definition of $\Gamma(\mathcal{L})$). We also choose distinct vertices $v_i \in V_{x_i}$ for $1 \leq i \leq r-1$ such that v_i is good for both triples \mathcal{H}_i and \mathcal{H}_{i+1} (here indices are taken mod $r-1$). If x_i is incident to the diamond arc (h, j) for some diamond $\{hjk, hjp\} \in \{\mathcal{D}_1, \dots, \mathcal{D}_s\}$, we also require that v_i be good for the two triples $\mathcal{H}[V_h, V_j, V_k]$ and $\mathcal{H}[V_h, V_j, V_p]$ as well. (Note that these two triples may or may not have been chosen as some \mathcal{H}_b .) Then Lemma 4.2 together with the fact that $r \leq 2t < \epsilon m$ (given by Lemma 2.2(3)) ensure that such a choice of $v_1, \dots, v_{r-1}, v_r = v_1$ can always be made. These vertices will be the end-vertices of our short paths.

Next we join v_{i-1} to v_i by a path of length three (or six) for each i , starting with $i = 2$. We apply Lemma 4.3 to the triple $\mathcal{H}_2 =$

$\mathcal{H}[V_{x_1}, V_{x_2}, V_{k_2}]$, the vertices $x = v_1$ and $y = v_2$, and the set B of vertices that are not good for \mathcal{H}_2 . This gives a path of length three joining v_1 and v_2 , whose two vertices of degree two lie in V_{x_1} and V_{x_2} .

In the general step we apply Lemma 4.3 to the triple $\mathcal{H}_i = \mathcal{H}[V_{x_{i-1}}, V_{x_i}, V_{k_i}]$, the vertices $x = v_{i-1}$ and $y = v_i$, and the set B of vertices that are not good for \mathcal{H}_i together with all vertices of previously defined short paths. Since the latter set has size at most $6r < \epsilon m$ by Lemma 2.2(3), the assumptions on the size of B in Lemma 4.3 are satisfied. This gives a path of length three joining v_{i-1} and v_i , disjoint from all previously defined paths, whose two vertices of degree two lie in $V_{x_{i-1}}$ and V_{x_i} . These short paths link to form a loose cycle of length $3(r-1)$. If the length of our target cycle \mathcal{C}_n has different parity from $3(r-1)$, we replace one path of length three that does not correspond to a diamond arc by a path of length six, using Corollary 4.4, to make the parities agree. This gives a loose cycle \mathcal{C} of length c , where $3(r-1) \leq c \leq 3r$ and $c \equiv n/2 \pmod{2}$.

Finally, to complete the proof we replace each short path in \mathcal{C} corresponding to a diamond arc by a long path. Let (x_{i-1}, x_i) be a diamond arc, for some diamond $\{x_{i-1}x_i k, x_{i-1}x_i p\}$. Then we apply Lemma 4.5 to the two ϵ -regular triples $\mathcal{H}[V_{x_{i-1}}, V_{x_i}, V_k]$ and $\mathcal{H}[V_{x_{i-1}}, V_{x_i}, V_p]$ of density $d \geq 1/2$, corresponding to the diamond. Here we are taking $x = v_{i-1}$ and $y = v_i$, and the length $\ell = \lfloor \frac{1}{s}(\frac{n}{2} - c) \rfloor$ or $\lfloor \frac{1}{s}(\frac{n}{2} - c) \rfloor - 1$, whichever is odd. We call this number the *basic length* for our long paths. Let the set B be the set of all vertices that are not good for $\mathcal{H}[V_{x_{i-1}}, V_{x_i}, V_k]$ or for $\mathcal{H}[V_{x_{i-1}}, V_{x_i}, V_p]$ together with the set of all vertices in previously defined paths. Note that the latter set still has size at most $6r < \epsilon m$, since the s diamonds $\mathcal{D}_1, \dots, \mathcal{D}_s$ of \mathcal{L} are vertex-disjoint and so the long paths cannot interfere with each other.

To check that the condition on ℓ is satisfied, recall from Lemma 2.3 that $s \geq f(\epsilon) \frac{n}{N} \frac{t(\epsilon_0)}{4}$, where $f(\epsilon) = (1 - 500\epsilon^{1/6})^{-2}$. From Theorem 2.1 we know $t \leq t(\epsilon_0)$ and $m \geq (1 - \epsilon) \frac{N}{t} > (1 - 500\epsilon^{1/6}) \frac{N}{t}$, from which it follows that

$$\frac{1}{s} \frac{n}{2} < (1 - 500\epsilon^{1/6}) 2m < (1 - 2\epsilon^{1/6}) 2m - 2r.$$

Therefore certainly $\lfloor \frac{1}{s}(\frac{n}{2} - c) \rfloor < (1 - 2\epsilon^{1/6}) 2m$ as required. Of course we also have $s \leq t$, so since $c \leq 3r \leq 6t < \epsilon m < \epsilon n$ (Lemma 2.2(3)) we know $\ell > \frac{1}{s}(\frac{n}{2} - c) - 2r > \frac{n}{t}(\frac{1}{2} - \epsilon) - 2r$. But then since $N < \frac{5}{4}n(1 + \eta) < \frac{3}{2}n$ (see Step 1), and by Theorem 2.1(2), this is at least $\frac{2}{3} \frac{N}{t}(\frac{1}{2} - \epsilon) - \epsilon m > \frac{2m}{7}$. Thus ℓ falls into the correct range for Lemma 4.5.

Therefore the assumptions of Lemma 4.5 are satisfied, and we can find a path joining v_i and v_{i+1} of length ℓ . For the last diamond we choose ℓ appropriately so that the resulting cycle has length exactly $n/2$, i.e. it is a copy of \mathcal{C}_n . This is possible since we would need to adjust the basic length by at most $2r$, and as above

$$\frac{2m}{7} < \frac{1}{s} \left(\frac{n}{2} - c \right) - 2r < \frac{1}{s} \left(\frac{n}{2} - c \right) + 2r \leq (1 - 2\epsilon^{1/6})2m.$$

This completes the proof. \square

5. PROOF OF LEMMA 2.3

We begin this section by focusing on a subhypergraph of \mathcal{J}_0 with convenient properties.

Lemma 5.1. *The hypergraph \mathcal{J}_0 contains a subhypergraph \mathcal{J}_1 with the following properties. Here $\epsilon_1 = 10\epsilon^{1/6}$.*

- (1) \mathcal{J}_1 has $t_1 > (1 - \epsilon_1)t$ vertices,
- (2) if some edge of \mathcal{J}_1 contains the vertices x and y then more than $(1 - \epsilon_1)t$ edges contain both x and y ,
- (3) for every vertex x , more than $(1 - \epsilon_1)t$ pairs xy are such that at least $(1 - \epsilon_1)t$ edges contain both x and y .

Proof. Recall from Step 1 (Lemma 2.2) that \mathcal{J}_0 has t vertices and at least $\binom{t}{3} - \epsilon t^3$ edges. We call a vertex x of \mathcal{J}_0 *bad* if it is in fewer than $\binom{t-1}{2} - \epsilon^{1/2}t^2$ edges. Then the number of bad vertices must be at most $3\epsilon^{1/2}t$. Remove these, and denote the resulting hypergraph by \mathcal{J}' . Then every vertex of \mathcal{J}' is in at least $\binom{t-1}{2} - \epsilon^{1/2}t^2 - 3\epsilon^{1/2}t^2 = \binom{t-1}{2} - 4\epsilon^{1/2}t^2$ edges.

We call a pair xy of vertices of \mathcal{J}' *bad* if it is contained in fewer than $(1 - 2\epsilon^{1/4})t$ edges. Then the maximum degree of the graph on $V(\mathcal{J}')$ formed by bad pairs is less than $4\epsilon^{1/4}t$, by the vertex property of \mathcal{J}' . Thus, in particular, the number of bad pairs is less than $4\epsilon^{1/4}t^2$.

We let \mathcal{J}_1 be the hypergraph formed by removing all edges that contain bad pairs. To check (1), note that $|V(\mathcal{J}_1)| = |V(\mathcal{J}')| > t - 3\epsilon^{1/2}t > (1 - \epsilon_1)t$.

Property (2) is true because if an edge remains, no pair in it was bad. Therefore if x and y are contained in this edge then they were contained in at least $(1 - 2\epsilon^{1/4})t$ edges of \mathcal{J}' . Thus, the number of edges lost for xy is at most $4\epsilon^{1/4}t$ for the bad pairs incident to x , plus $4\epsilon^{1/4}t$ for the bad pairs incident to y . Therefore xy is contained in at least $(1 - 10\epsilon^{1/4})t > (1 - \epsilon_1)t$ edges of \mathcal{J}_1 .

Finally to verify (3), note that the above argument tells us that each pair xy that was not bad satisfies the property in (2). Thus x is incident

to more than $|V(\mathcal{J}_1)| - 4\epsilon^{1/4}t > (1 - 3\epsilon^{1/2} - 4\epsilon^{1/4})t > (1 - \epsilon_1)t$ such pairs.

□

The following lemma is a simple calculation that relates s to the parameters of \mathcal{J}_1 .

Lemma 5.2. *We have $s \leq \frac{1}{5}|V(\mathcal{J}_1)| - 6\epsilon_1t$.*

Proof. Recall that $s = \lceil f(\epsilon)\frac{n}{N}\frac{t}{4} \rceil$, where $f(\epsilon)$ is as defined in Lemma 2.3, and $N = 5g(\epsilon_0)n/4$ (see Step 1). Then since $\epsilon = \epsilon_0^{1/4}$ (see Lemma 2.2) and $\epsilon_1 = 10\epsilon^{1/6}$ (see Lemma 5.1) we have $N = 5(1 - 50\epsilon_1)^{-1}f(\epsilon)n/4$. Therefore

$$f(\epsilon)\frac{n}{N}\frac{5t}{4} = (1 - 50\epsilon_1)t = (1 - \epsilon_1)t - 49\epsilon_1t.$$

Hence $f(\epsilon)\frac{n}{N}\frac{t}{4} < \frac{1}{5}(1 - \epsilon_1)t - 9\epsilon_1t$, which implies the desired result because $|V(\mathcal{J}_1)| > (1 - \epsilon_1)t$ by Lemma 5.1(1). □

We are interested in the monochromatic components of the 2-coloured hypergraph \mathcal{J}_1 (as defined in Step 2). Note that each vertex of \mathcal{J}_1 is in one red monochromatic component and one blue monochromatic component.

Lemma 5.3. *The hypergraph \mathcal{J}_1 has the following properties.*

- (1) \mathcal{J}_1 contains a monochromatic component with more than $(1 - 3\epsilon_1)t$ vertices. Let us say without loss of generality that this component \mathcal{R} is red.
- (2) if \mathcal{R} does not contain s disjoint diamonds, then the largest blue component \mathcal{B} has at least $4|V(\mathcal{J}_1)|/5$ vertices.

Proof. For the first assertion, consider the shadow graph $G = \Gamma(\mathcal{J}_1)$ as defined in Step 2. Then the minimum degree of G is more than $(1 - \epsilon_1)t$ by Lemma 5.1(3). We call an edge xy of G *red* if it is contained in a red edge of \mathcal{J}_1 , and *blue* if it is in a blue edge of \mathcal{J}_1 . Note that some edges may be both red and blue. Then a monochromatic component of \mathcal{J}_1 is by definition the same as a monochromatic component of G . Suppose that the largest monochromatic (say blue) component C in G does not cover all vertices of G . Certainly $|C| > (1 - \epsilon_1)t/2$. Let x be a vertex outside C . Then all the edges of G joining x to a vertex of C are red, and there are more than $|C| - \epsilon_1t$ such edges. Let $y \in C$ be a neighbour of x , then all the edges of G joining y to $V(G) \setminus C$ are red as well, and there are more than $|V(G) \setminus C| - \epsilon_1t > (1 - 2\epsilon_1)t - |C|$ such edges. Then this gives a red component in G of size more than $(1 - 3\epsilon_1)t$ as required.

To prove (2), suppose the largest blue component \mathcal{B} has at most $4|V(G)|/5$ vertices. We first claim that there exists a set $A \subset V(G)$ with $|V(G)|/5 \leq |A| \leq |V(G)|/2$ such that all edges of G that contain a vertex of A and a vertex of $\bar{A} = V(G) \setminus A$ are red but not blue. This is clearly true if $|V(\mathcal{B})| \geq |V(G)|/5$, since either $V(\mathcal{B})$ or $V(G) \setminus V(\mathcal{B})$ would do. If $|V(\mathcal{B})| < |V(G)|/5$ then there is a subset S of the blue components such that $A' = \bigcup_{\mathcal{B}_i \in S} V(\mathcal{B}_i)$ satisfies $2|V(G)|/5 \leq |A'| \leq 3|V(G)|/5$. (Note that some blue components may just be isolated vertices.) Then again either A' or $V(\mathcal{J}_1) \setminus A'$ is a suitable choice for A .

Note then that the red component of \mathcal{J}_1 given by these red edges of G must be \mathcal{R} . Next we show that this structure for \mathcal{R} guarantees that it contains s disjoint diamonds. We construct greedily a set of disjoint diamonds in \mathcal{R} as follows. Suppose $i < \frac{1}{6}(4|A| - |\bar{A}|) - \epsilon_1 t$, and that disjoint diamonds $\mathcal{D}_1, \dots, \mathcal{D}_{i-1}$ have already been found in \mathcal{R} , each of which has two vertices in each of A and \bar{A} . Then there remain at least $|A| - 2(i-1) \geq 2\epsilon_1 t$ unused vertices in A (and hence also in \bar{A}). Then by Lemma 5.1(3), there is an edge xy of G with $x \in A$ and $y \in \bar{A}$ both unused, which is contained in more than $(1 - \epsilon_1)t$ edges of \mathcal{J}_1 , which are all red. Therefore there exists another unused vertex $w \in A$, and another unused $z \in \bar{A}$, such that xyw and xyz are both red edges of \mathcal{J}_1 . Then we set \mathcal{D}_i to be this red diamond.

When we have found $\lceil \frac{1}{6}(4|A| - |\bar{A}|) - \epsilon_1 t \rceil$ diamonds in \mathcal{R} as above, we add some more diamonds that have one vertex in A and three in \bar{A} . Suppose $i < \frac{1}{3}(|\bar{A}| - |A|)$, and that disjoint diamonds $\mathcal{D}_1, \dots, \mathcal{D}_{i-1}$ have already been found in \mathcal{R} , each of which has one vertex in A and three in \bar{A} . Then there remain at least $|A| - 2\lceil \frac{1}{6}(4|A| - |\bar{A}|) - \epsilon_1 t \rceil - i - 1 \geq 2\epsilon_1 t$ unused vertices in A and $|\bar{A}| - 2\lceil \frac{1}{6}(4|A| - |\bar{A}|) - \epsilon_1 t \rceil - 3(i-1) \geq 2\epsilon_1 t$ in \bar{A} . Then by Lemma 5.1(3), there is an edge xy of G with $x \in A$ and $y \in \bar{A}$ both unused, which is contained in more than $(1 - \epsilon_1)t$ edges of \mathcal{J}_1 , which are all red. Therefore there exist two more unused vertices $w, z \in \bar{A}$, such that xyw and xyz are both red edges of \mathcal{J}_1 . Then we set \mathcal{D}_i to be this red diamond.

The total number of disjoint diamonds we find by this greedy construction is at least $\frac{1}{6}(4|A| - |\bar{A}|) - \epsilon_1 t + \frac{1}{3}(|\bar{A}| - |A|) \geq \frac{|V(G)|}{5} - \epsilon_1 t$, which is more than s by Lemma 5.2. \square

In addition to Lemma 5.3, the following technical lemma will be important for the proof of Lemma 2.3. By a *diadem* we mean any coloured hypergraph with 5 vertices that contains both a red diamond and a blue diamond.

Lemma 5.4. *Let $X = \{x_1, x_2, x_3, x_4\}$ and $Y = \{y_1, y_2, y_3, y_4, y_5\}$ be disjoint vertex sets such that $x_i x_j y_k \in \mathcal{J}_1$ for all i, j, k except possibly when $\{i, j\} = \{1, 4\}$. Let $S \subset X$ be a subset of size at least two. Suppose further that $x_1 x_2 x_3, x_2 x_3 x_4 \in \mathcal{J}_1$ and are both red. Then one of the following holds.*

- (1) $X \cup Y$ contains two disjoint red diamonds, both of which intersect X .
- (2) $X \cup Y$ contains a diadem \mathcal{E} such that $3 \leq |V(\mathcal{E}) \cap X| \leq 4$ whose blue diamond intersects both S and Y .

Proof. Since $|S| \geq 2$, either we have both $\{x_1, x_2\} \cap S \neq \emptyset$ and $\{x_3, x_4\} \cap S \neq \emptyset$, or $\{x_1, x_3\} \cap S \neq \emptyset$ and $\{x_2, x_4\} \cap S \neq \emptyset$. Without loss of generality we may assume that the former possibility holds. We look first at the edges $x_1 x_2 y_i$, $i = 1, \dots, 5$. We argue that if more than one of them is blue, then $X \cup Y$ contains a diadem.

Indeed, suppose that, say, the edges $x_1 x_2 y_1$ and $x_1 x_2 y_2$ are blue. Then, if the edge $x_1 x_3 y_1$ is blue, we take the vertex set W of \mathcal{E} to be $W = \{x_1, x_2, x_3, x_4, y_1\}$; if $x_1 x_3 y_1$ is red, then we put $W = \{x_1, x_2, x_3, y_1, y_2\}$. Consequently, we may assume that, say, the edges $x_1 x_2 y_i$, $i = 1, \dots, 4$, are coloured red.

An analogous argument shows that, to avoid a diadem, at most one of the edges $x_3 x_4 y_i$, $i = 1, \dots, 4$, is blue, so suppose that $x_3 x_4 y_i$, $i = 1, 2, 3$, are red. But then $X \cup Y$ contains two red diamonds with edges $x_1 x_2 y_4$, $x_1 x_2 y_3$, $x_3 x_4 y_2$ and $x_3 x_4 y_1$. \square

We are now ready to prove Lemma 2.3.

Proof of Lemma 2.3. Let $\{\mathcal{D}_1, \dots, \mathcal{D}_q\}$ be a set of disjoint diamonds in \mathcal{R} of maximum size q . If $q \geq s$ we have proved the lemma, so let us assume $q < s$. Our aim is to show that the largest blue component \mathcal{B} contains s disjoint diamonds.

Our first step is to replace the red diamonds in \mathcal{R} one-by-one by disjoint diadems. Suppose $0 \leq i < q$, let $Q_i = \bigcup_{j=i+1}^q V(\mathcal{D}_j)$, and suppose that disjoint diadems $\mathcal{E}_1, \dots, \mathcal{E}_i$ have been found such that

- $V(\mathcal{E}_1 \cup \dots \cup \mathcal{E}_i) \cap Q_i = \emptyset$,
- each \mathcal{E}_j has its red diamond in \mathcal{R} and its blue diamond in \mathcal{B} ,
- the set $V_i = V(\mathcal{J}_1) \setminus V(\mathcal{E}_1 \cup \dots \cup \mathcal{E}_i)$ contains at least $\frac{4}{5}|V_i|$ vertices of \mathcal{B} .

Lemma 5.3(2) guarantees that this process can start.

We now consider three cases.

Case 1. There exists $j \in \{i+1, \dots, q\}$, say without loss of generality $j = i+1$, such that \mathcal{D}_j has at least 2 vertices in \mathcal{B} and $V_i \setminus Q_j$ has at least $5\epsilon_1 t + 5$ vertices not in \mathcal{B} .

Set $X = V(\mathcal{D}_j)$. By Lemma 5.1(2) there exists $Y \subset V_i \setminus (V(\mathcal{B}) \cup Q_i)$, $|Y| = 5$, such that $X \cup Y$ satisfies the conditions of Lemma 5.4. We let $S = X \cap V(\mathcal{B})$ and apply Lemma 5.4. Now note that outcome (1) is not possible, since otherwise we would have a set of disjoint red diamonds of size $q+1$ in \mathcal{R} , contradicting the definition of q . Therefore we have a diadem \mathcal{E} satisfying (2). But then the blue diamond in \mathcal{E} intersects S and hence is in \mathcal{B} , the red diamond of \mathcal{E} intersects X and hence is in \mathcal{R} , and $|V(\mathcal{E}) \cap V(\mathcal{B})| \leq 4$ because $Y \cap V(\mathcal{B}) = \emptyset$. Therefore $V_{i+1} = V_i \setminus V(\mathcal{E})$ contains at least $\frac{4}{5}|V_{i+1}|$ vertices of \mathcal{B} . Thus setting $\mathcal{E}_{i+1} = \mathcal{E}$ allows us to continue the construction.

Case 2. $|(V_i \setminus Q_i) \cap V(\mathcal{B})| \geq |V_i \setminus Q_i| - 5\epsilon_1 t - 5$.

Note that $|V_i \setminus Q_i| \geq |V(\mathcal{J}_1)| - 5q > |V(\mathcal{J}_1)| - 5s \geq 10\epsilon_1 t + 10$ by Lemma 5.2, so $|(V_i \setminus Q_i) \cap V(\mathcal{B})| \geq 5\epsilon_1 t + 5$. We choose $j \in \{i+1, \dots, q\}$ such that \mathcal{D}_j has the smallest possible number p of vertices of \mathcal{B} . Then with $X = V(\mathcal{D}_j)$, as before there exists $Y \subset (V_i \setminus Q_i) \cap V(\mathcal{B})$, $|Y| = 5$, such that $X \cup Y$ satisfies the conditions of Lemma 5.4. Again we can't have two disjoint red diamonds, so there is a diadem \mathcal{E} whose blue diamond intersects Y . This guarantees that the blue diamond of \mathcal{E} is in \mathcal{B} . Since $|V(\mathcal{E}) \cap X| \geq 3$ the red diamond of \mathcal{E} is in \mathcal{R} .

Now if $p \leq 2$ then $|V(\mathcal{E}) \cap V(\mathcal{B})| \leq 4$ and we find that $V_{i+1} = V_i \setminus V(\mathcal{E})$ contains at least $\frac{4}{5}|V_{i+1}|$ vertices of \mathcal{B} as in Case 1. But if $p \geq 3$ then by definition of p we know

$$\begin{aligned} |V_i \cap V(\mathcal{B})| &\geq \frac{3}{4}|V_i \cap Q_i| + |(V_i \setminus Q_i) \cap V(\mathcal{B})| \\ &\geq 3q - 3i + |V_i \setminus Q_i| - 5\epsilon_1 t - 5 \\ &\geq 3q - 3i + |V(\mathcal{J}_1)| - 4q - i - 5\epsilon_1 t - 5 \\ &= |V(\mathcal{J}_1)| - q - 4i - 5\epsilon_1 t - 5 \\ &= \left(\frac{4}{5}|V(\mathcal{J}_1)| - 4i\right) + \frac{1}{5}|V(\mathcal{J}_1)| - q - 5\epsilon_1 t - 5 \\ &\geq \frac{4}{5}|V_i| + 1, \end{aligned}$$

where the last line follows since $|V_i| = |V(\mathcal{J}_1)| - 5i$ and $q < s \leq \frac{1}{5}|V(\mathcal{J}_1)| - 5\epsilon_1 t - 6$ by Lemma 5.2. Therefore $|V_{i+1} \cap V(\mathcal{B})| \geq \frac{4}{5}|V_i| + 4 = \frac{4}{5}(|V_i| - 5) = \frac{4}{5}|V_{i+1}|$ as required.

Case 3. Each \mathcal{D}_j with $j \in \{i+1, \dots, q\}$ has at most one vertex of \mathcal{B} .

Now if $|(V_i \setminus Q_i) \cap V(\mathcal{B})| \geq 5\epsilon_1 t + 5$ then as in Case 2 we can find a diadem \mathcal{E} with its blue diamond in \mathcal{B} , its red diamond in \mathcal{R} , and $|V(\mathcal{E}) \cap V(\mathcal{B})| \leq 3$. Then $\mathcal{E}_{i+1} = \mathcal{E}$ is a suitable choice. But if $|(V_i \setminus$

$|Q_i \cap V(\mathcal{B})| < 5\epsilon_1 t + 5$ then

$$\begin{aligned} \frac{4}{5}|V(\mathcal{J}_1)| - 4i &= \frac{4}{5}|V_i| \leq |V_i \cap V(\mathcal{B})| = |V_i \cap Q_i \cap V(\mathcal{B})| + |(V_i \setminus Q_i) \cap V(\mathcal{B})| \\ &\leq q - i + 5\epsilon_1 t + 5. \end{aligned}$$

Therefore $\frac{4}{5}|V(\mathcal{J}_1)| \leq q + 3i + 5\epsilon_1 t + 5 < 4s + 5\epsilon_1 t + 5$, implying $\frac{1}{5}|V(\mathcal{J}_1)| < s + \frac{5}{4}\epsilon_1 t + \frac{5}{4}$, which contradicts Lemma 5.2. Thus this cannot occur.

Therefore we can complete the above construction to find disjoint diadems $\mathcal{E}_1, \dots, \mathcal{E}_q$, whose red diamonds are all in \mathcal{R} and whose blue diamonds are all in \mathcal{B} , such that $V_q = V(\mathcal{J}_1) \setminus V(\mathcal{E}_1 \cup \dots \cup \mathcal{E}_q)$ contains at least $\frac{4}{5}|V_q|$ vertices of \mathcal{B} . Note that $|V_q| = |V(\mathcal{J}_1)| - 5q \geq 30\epsilon_1 t$ by Lemma 5.2. But by definition of q , no red diamond can be present in $V_q \cap V(\mathcal{R})$. Therefore for any $x \in V_q \cap V(\mathcal{R})$ and any $y \in V_q \cap V(\mathcal{R})$ such that xy is in at least $(1 - \epsilon_1)t$ edges of \mathcal{J}_1 (which exists by Lemma 5.1(3) since $|V_q \cap V(\mathcal{R})| > 2\epsilon_1 t$), all but at most $\epsilon_1 t + 1$ of the vertices $z \in V_q \cap V(\mathcal{R})$ are such that $xyz \in \mathcal{J}_1$ and xyz is blue. Therefore in particular

$$\begin{aligned} |V_q \cap V(\mathcal{R}) \cap V(\mathcal{B})| &\geq |V_q \cap V(\mathcal{R})| - \epsilon_1 t - 1 \\ &\geq |V(\mathcal{J}_1)| - 5q - 4\epsilon_1 t - 1 > 4(s - q) + 2\epsilon_1 t, \end{aligned}$$

where the last line follows since $|V(\mathcal{J}_1)| > 5s + 30\epsilon_1 t > 4s + q + 6\epsilon_1 t + 1$ by Lemma 5.2. Thus we can easily find a set U of $s - q$ blue diamonds in $V_q \cap V(\mathcal{R}) \cap V(\mathcal{B})$, just by a greedy construction. Then U together with the q blue diamonds in $\mathcal{E}_1, \dots, \mathcal{E}_q$ give a set of s disjoint blue diamonds in \mathcal{B} as required. \square

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