SHORT PATHS IN QUASI-RANDOM TRIPLE SYSTEMS WITH SPARSE UNDERLYING GRAPHS

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Abstract. The regularity lemma for 3-uniform hypergraphs asserts that every large hypergraph can be decomposed into a bounded number of quasi-random structures consisting of a sub-hypergraph and a sparse underlying graph. In this paper we show that in such a quasi-random structure most pairs of the edges of the graph can be connected by hyperpaths of length at most twelve. Some applications are also given.

1. Introduction

The Regularity Lemma from [10] is a powerful tool in contemporary graph theory and combinatorics. It allows one to partition every large graph into a bounded number of bipartite subgraphs, most of which are quasi-random, that is, they possess essentially all typical properties of corresponding random graphs. One of these properties, quite easy to prove, is that every two vertices with non-negligible neighborhoods can be connected by a path of length at most four (see, e.g., [7] and Corollary 2.5(a)).

In this paper we study the much harder problem of the existence of short paths in 3-uniform, 3-partite hypergraphs with a certain regular structure related to the Hypergraph Regularity Lemma in [2]. When this lemma is being applied, the initial hypergraph is broken into several quasi-random pieces and a desired structure is built from segments scattered among these highly regular substructures. It is then important to “sew” them together by relatively short hyperpaths.

Two examples of this general approach can be found in the forthcoming papers [9] and [4], where, respectively, the existence of Hamilton cycles in 3-uniform hypergraphs and the Ramsey numbers for hypercycles are treated. In each of these applications, besides the Hypergraph Regularity Lemma itself, a crucial role is played by a “connection lemma” guaranteeing short paths between (almost) all pairs of pairs of vertices.

In [9] such a lemma follows from the strong assumption that every pair of vertices is contained in more than $n/2$ hyperedges. The connection lemma applied in [4] is, on the other hand, a consequence of the quasi-random structure and as such is analogous to, but much more complicated than its counterpart for graphs. In Section 7.1, we describe this application in more detail.
The goal of this paper is to prove the connection lemma for quasi-random, 3-uniform hypergraphs, in the form stated in [4]. In the next section, after some preliminary definitions, we state our main result, Theorem 2.16. Then, in Section 3 we reformulate it in a more constructive way, specifying, in terms of their fourth neighborhoods, the edges that can be connected by short hyperpaths. Section 4 contains the proofs of these two theorems, both relying on two lemmas, Lemma 4.1 and Lemma 4.2, which themselves will be proved in Sections 5 and 6. Section 7 presents briefly some applications of Theorem 2.16. One of them guarantees a sub-hamiltonian path in a quasi-random 3-uniform hypergraph and, in turn, is used to derive asymptotic values of the Ramsey numbers for hypercycles in [4]. In the final application, we approximate every large 3-uniform hypergraph by finitely many pieces of small “diameter”.

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2. Preliminaries and main result

2.1. Facts on \(\epsilon\)-regular pairs. In this subsection we collect elementary facts about \(\epsilon\)-regular graphs which are used throughout the paper.

Let \(G = (V, E)\) be a graph, where \(V\) and \(E\) are the vertex-set and the edge-set of \(G\). Throughout the paper we often identify \(G\) with its set of edges and therefore write \(|G|\) instead of \(|E|\). When \(U\) and \(W\) are subsets of \(V\), we define

\[ e_G(U, W) = |\{x, y \in E : x \in U, y \in W\}|. \]

For nonempty and disjoint \(U\) and \(W\),

\[ d_G(U, W) = \frac{e_G(U, W)}{|U||W|} \]

is the density of the graph \(G\) between \(U\) and \(W\), or simply, the density of the pair \((U, W)\).

**Definition 2.1.** Given \(\epsilon > 0\), a bipartite graph \(G\) with bipartition \((V_1, V_2)\), where \(|V_1| = n\) and \(|V_2| = m\), is called \(\epsilon\)-regular if for every pair of subsets \(U \subseteq V_1\) and \(W \subseteq V_2\), \(|U| > \epsilon n\), \(|W| > \epsilon m\), the inequalities

\[ d - \epsilon < d_G(U, W) < d + \epsilon \]

hold for some real number \(d > 0\). We may then also say that \(G\), or the pair \((V_1, V_2)\), is \((d, \epsilon)\)-regular.

Let a graph \(G = (V, E)\) be given. We write \(N_G(v)\) for the set of neighbors of \(v \in V\) in the graph \(G\). The size of \(N_G(v)\) is \(|N_G(v)| = \deg_G(v)\), the degree of \(v\). We set \(N_G(xy) = N_G(x) \cap N_G(y)\) as the set of common neighbors of \(x, y \in V\) in \(G\). For a set \(U \subseteq V\), we write \(N_G(v, U)\) for the set of neighbors of \(v\) in \(U\) and \(N_G(xy, U)\) for the set of common neighbors of \(x\) and \(y\) in \(U\). The size of \(N_G(v, U)\) is \(|N_G(v, U)| = \deg_G(v, U)|\).
Definition 2.2. Let $G = (V_1 \cup V_2, E)$ be a $(d, \epsilon)$-regular bipartite graph, where $|V_1| = |V_2| = n$. We say that a vertex $x \in V_i$, $i = 1, 2$, is typical in $G$, if the following inequalities hold

\[ n(d - \epsilon) < \deg_G(x) < (d + \epsilon)n. \]

Further, let $G = G^{12} \cup G^{23} \cup G^{13}$ be a 3-partite graph with partition $(V_1, V_2, V_3)$, where $|V_1| = |V_2| = |V_3| = n$, and each graph $G^{ij}$ is $(d, \epsilon)$-regular, $1 \leq i < j \leq 3$. We call a pair of vertices $(x, y) \in V_i \times V_j$ typical if it satisfies inequalities

\[ n(d - \epsilon)^2 < |N_G(xy)| < n(d + \epsilon)^2. \]

The next fact is well-known and follows immediately from Definition 2.1 (see e.g. [1],[7]).

Fact 2.3. For all $\epsilon > 0$ and $d > 0$, and for all integers $n$ and $m$, the following holds. Let $G$ be a $(d, \epsilon)$-regular bipartite graph with a bipartition $(V_1, V_2)$, where $|V_1| = n, |V_2| = m$. Further, let $A \subseteq V_2, |A| > em$. Then all but at most $en$ vertices $x \in V_1$ satisfy

\[ \text{deg}_G(x, A) < (d + \epsilon)|A|, \]

and all but at most $en$ vertices $x \in V_1$ satisfy

\[ \text{deg}_G(x, A) > (d - \epsilon)|A|. \]

In particular, if $|V_1| = |V_2| = n$, then for each $i \in \{1, 2\}$, all but at most $2en$ vertices $x \in V_i$ are typical in $G$.

Corollary 2.4. For all $\epsilon > 0$ and $d > 2\epsilon$ and for all integers $n$, the following holds. Let $G = G^{12} \cup G^{23} \cup G^{13}$ be a 3-partite graph with partition $(V_1, V_2, V_3)$, where $|V_1| = |V_2| = |V_3| = n$ and each graph $G^{ij}$ is $(d, \epsilon)$-regular, $1 \leq i < j \leq 3$. Then all but at most $4en^2$ pairs of vertices $(x, y) \in V_i \times V_j$ are typical.

Another simple consequence of Fact 2.3 deals with the distances in a quasi-random bipartite graph (see [7] and [8]).

Corollary 2.5. Let $B$ be a $(d, \epsilon)$-regular bipartite graph with bipartition $(V_1, V_2)$, where $|V_1| = |V_2| = n$.

(a) If $d > 2\epsilon$ then all pairs of vertices of $B$ of degree at least $en$ can be connected by paths of length at most four.

(b) If $d > 4\epsilon$ then by removing from $B$ at most $2en$ vertices (those of degree less than $3en < (d - \epsilon)n$), we obtain a subgraph with diameter four.

Finally, we state another well-known result which tightly estimates the size of $\text{Tr}(G)$, the set of triangles in a quasi-random 3-partite graph $G$ (see, e.g., [2],[9]).

Fact 2.6. Let $G = G^{12} \cup G^{23} \cup G^{13}$ be a 3-partite graph, where all three bipartite graphs $G^{ij}$ are $(d, \epsilon)$-regular, $1 \leq i < j \leq 3$. If $d > 2\epsilon$ then

\[ (d^3 - 10\epsilon) < \frac{|\text{Tr}(G)|}{|V_1||V_2||V_3|} < (d^3 + 10\epsilon). \]

In particular, if $\epsilon < 0.1d^3$ then $|\text{Tr}(G)| < 2d^3|V_1||V_2||V_3|$. 

\[ \frac{|\text{Tr}(G)|}{|V_1||V_2||V_3|} < (d^3 + 10\epsilon). \]
2.2. Regularity of hypergraphs. We begin with some basic definitions from hypergraph theory.

Definition 2.7. A 3-uniform hypergraph is a pair $\mathcal{H} = (V, E)$, where $V$ is a finite set of vertices and $E$ is a family of 3-element subsets of $V$ called hyperedges or triplets. Throughout the paper we will often identify $\mathcal{H}$ with $E$.

We call $\mathcal{H}$ 3-partite if there exists a partition $V = V_1 \cup V_2 \cup V_3$ such that for each $e \in E$ and for each $i = 1, 2, 3$ we have $e \cap V_i \neq \emptyset$. We refer to any 3-partite 3-uniform hypergraph $\mathcal{H}$ with a fixed 3-partition $(V_1, V_2, V_3)$ as a 3-graph.

For an arbitrary hypergraph $\mathcal{H}$ and a graph $G$ on the same vertex set, we denote by $\mathcal{H} - G$ the sub-hypergraph of $\mathcal{H}$ obtained by removing all hyperedges containing at least one edge of $G$.

The density and $\epsilon$-regularity of bipartite graphs is measured by the ratio of the number of edges to all potential edges (see above). For 3-graphs it is the ratio of hyperedges coinciding with the triangles of an underlying graph to all triangles in that graph.

Definition 2.8. For a 3-partite graph $P$ with a fixed 3-partition $V_1 \cup V_2 \cup V_3$, we shall write $P = P^{12} \cup P^{23} \cup P^{13}$, where $P^{ij} = \{xy \in P : x \in V_i, y \in V_j\}$. Furthermore, let $\text{Tr}(P)$ be the set of all (vertex sets of) triangles formed by the edges of $P$. If $P = P^{12} \cup P^{23} \cup P^{13}$ is a 3-partite graph with the same vertex partition as $\mathcal{H}$, and moreover, $\mathcal{H} \subseteq \text{Tr}(P)$, then we say that $P$ underlies $\mathcal{H}$.

The natural notion of density $d_{\mathcal{H}}(P)$ of $\mathcal{H}$ with respect to $P$ counts the proportion of triangles of $P$ which are triplets of $\mathcal{H}$. Then, the $\delta$-regularity of $\mathcal{H}$ means that for all $Q \subseteq P$ that contain at least $\delta |\text{Tr}(P)|$ triangles, the densities of $\mathcal{H}$ with respect to such $Q$'s are within $\delta$ from each other. However, it turns out that in some applications this is not strong enough. Therefore, the concept of so called $(\delta, r)$-regularity was introduced in [2].

Definition 2.9. Let $r \geq 1$ be an integer and let $\mathcal{H}$ be a 3-graph with an underlying 3-partite graph $P = P^{12} \cup P^{23} \cup P^{13}$. Let $Q = (Q(1), \ldots, Q(r))$ be an $r$-tuple of 3-partite subgraphs $Q(s) = Q^{12}(s) \cup Q^{23}(s) \cup Q^{13}(s)$ satisfying that for all $s \in \{1, 2, \ldots, r\}$ and $1 \leq i < j \leq 3$, $Q^{ij}(s) \subseteq P^{ij}$. We define the density $d_{\mathcal{H}}(Q)$ of $\mathcal{H}$ with respect to $Q$ as

\[
d_{\mathcal{H}}(Q) = \frac{|\mathcal{H} \cap \bigcup_{s=1}^{r} \text{Tr}(Q(s))|}{|\bigcup_{s=1}^{r} \text{Tr}(Q(s))|},
\]

if $|\bigcup_{s=1}^{r} \text{Tr}(Q(s))| > 0$, and 0 otherwise.

Definition 2.10. Let an integer $r \geq 1$ and real numbers $0 < \alpha, \delta < 1$ be given. We say that a 3-graph $\mathcal{H}$ is $(\alpha, \delta, r)$-regular with respect to an underlying graph $P = P^{12} \cup P^{23} \cup P^{13}$ if for any $r$-tuple of subgraphs $Q = (Q(1), \ldots, Q(r))$ as above, if

\[
\left| \bigcup_{s=1}^{r} \text{Tr}(Q(s)) \right| > \delta |\text{Tr}(P)|,
\]

then

\[
|d_{\mathcal{H}}(Q) - \alpha| < \delta.
\]
We say that $\mathcal{H}$ is $(\delta, r)$-regular with respect to $P$ if it is $(\alpha, \delta, r)$-regular for $\alpha = d_{\mathcal{H}}(P)$. Note that if $\mathcal{H}$ is $(\delta, r)$-regular with respect to $P$, $\delta' \geq \delta$, and $r' \leq r$ is an integer, then $\mathcal{H}$ is also $(\delta', r')$-regular with respect to $P$. If $r = 1$, we just use the names $\delta$-regular and $(\alpha, \delta)$-regular.

**Setup 2.11.** In what follows we always assume that $\mathcal{H}$ is a 3-graph and $P = P_1^{12} \cup P_2^{23} \cup P_3^{13}$ is a 3-partite graph, both with the same 3-partition $V = V(\mathcal{H}) = V(P) = V_1 \cup V_2 \cup V_3$ with $|V_1| = |V_2| = |V_3| = n$, and moreover, that $P$ underlies $\mathcal{H}$, i.e., $\mathcal{H} \subseteq \text{Tr}(P)$.

**Definition 2.12.** Given $\mathcal{H}$ and $P$ as in Setup 2.11, integers $l$ and $r$ and real numbers $\alpha$, $\delta$ and $\epsilon$, we call the pair $(\mathcal{H}, P)$ an $(\alpha, \delta, l, r, \epsilon)$-triad if

1. each $P_{ij}^l$, $1 \leq i < j \leq 3$, is $(1/l, \epsilon)$-regular;
2. $\mathcal{H}$ is $(\alpha, \delta, r)$-regular with respect to $P$.

We call $(\mathcal{H}, P)$ an $(\geq \alpha, \delta, l, r, \epsilon)$-triad if is a $(\beta, \delta, l, r, \epsilon)$-triad for some $\beta \geq \alpha$. In particular, it follows that if $(\mathcal{H}, P)$ is an $(\alpha, \delta, l, r, \epsilon)$-triad then for all $1 \leq i < j \leq 3$ we have

$$
(1/l - \epsilon)n^2 < |P_{ij}| < (1/l + \epsilon)n^2.
$$

The Hypergraph Regularity Lemma in [2] states that with the right choice of parameters, for every large and dense 3-uniform hypergraph $\mathcal{H} = (V, E)$, the complete graph on $V$ can be partitioned into finitely many graphs so that most triplet s of $\mathcal{H}$ belong to $(\alpha, \delta, l, r, \epsilon)$-triads built upon these graphs. This paper studies the structure of $(\mathcal{H}, P)$ in such a typical situation.

For future references in Section 7, we now state the regularity lemma for 3-uniform hypergraphs from [2] in a simplified form presented in [9] (see Lemma 4.1 and Remark 4.1 there). Set $K(U, W)$ for the complete bipartite graph with vertex sets $U$ and $W$.

**Theorem 2.13** (The Hypergraph Regularity Lemma). For every $\delta > 0$, every integer $t_0$ and for all integer-valued functions $r = r(t, l)$ and all decreasing functions $\epsilon(l) > 0$, there exist constants $T_0, L_0$ and $N_0$ such that every 3-uniform hypergraph $\mathcal{H}$ with at least $N_0$ vertices admits a partition $\Pi$ consisting of an auxiliary vertex set partition $V(\mathcal{H}) = V_0 \cup V_1 \cup \cdots \cup V_r$ where $t_0 \leq t < T_0$, $|V_0| < t$ and $|V_1| = |V_2| = \cdots = |V_t|$, and, for each pair $i$, $j$, $1 \leq i < j \leq t$, of a partition $K(V_i, V_j) = \bigcup_{l=1}^{r} P_{ij}^l$, where $1 \leq l < L_0$, satisfying the following conditions:

1. all graphs $P_{ij}^l$ are $(1/l, \epsilon(l))$-regular;
2. $\mathcal{H}$ is $(\delta, r)$-regular with respect to all but at most $\delta^3 t^3$ triads $(P_{ij}^l, P_{ij}^b, P_{ij}^c)$.

2.3. Main result. There are several ways to define a path in a 3-uniform hypergraph, and we choose one in which the edges are glued along the path in the most tight way (see [6] and [3] for some study of paths and cycles defined in a “loose” way).

**Definition 2.14.** Let $\mathcal{H}$ be a 3-uniform hypergraph. A hyperpath of length $k \geq 0$ in $\mathcal{H}$ is a sub-hypergraph $\mathcal{P}$ of $\mathcal{H}$ consisting of $k + 2$ vertices and $k$ hyperedges and whose vertices can be labeled $x_0, \ldots, x_{k+2}$ so that for each $i = 1, \ldots, k$, $x_ix_{i+1}x_{i+2} \in \mathcal{H}$. We then say that $\mathcal{P}$ goes from the pair $x_1x_2$ to the pair $x_{k+2}x_{k+1}$ and these two pairs are called the endpairs of $\mathcal{P}$. The vertices $x_3, \ldots, x_k$ are called internal. Two paths are said to be internally disjoint if they do not share any internal vertex.
Remark 2.15. Note that the endpairs are ordered pairs of vertices. If $\mathcal{H}$ is a 3-partite hypergraph then the vertices of any hyperpath traverse the partition sets only in the cyclic order $V_1 \to V_2 \to V_3 \to V_1$, or in its reverse (see Figure 1). Hence, there are pairs of ordered pairs of vertices which, even in a complete 3-graph, are not connected by any hyperpath. Another consequence is that the lengths of paths connecting two given endpairs are equal modulo 3.

Throughout the paper we will be assuming that the cyclic ordering $V_1 \to V_2 \to V_3 \to V_1$ is canonical, and thus, specifying two unordered pairs of vertices, $e$ and $f$, and saying that a hyperpath goes from $e$ to $f$ will not be ambiguous. (Note that under this convention a hyperpath from $f$ to $e$ is not a mere reverse of a path from $e$ to $f$.)

Note also that unlike the graph case, the length of the shortest hyperpath between two given endpairs does not satisfy the triangle inequality, and thus cannot be called “distance”.

Our goal is to prove the following “Connection Lemma” which, in a way, extends a simple fact about graphs, Corollary 2.5(b) (see above), to 3-uniform quasi-random hypergraphs. In addition, for the sake of future applications, we may force the hyperpaths to avoid a specified set of vertices $S$. A hyperpath $\mathcal{P}$ is called $S$-avoiding if $V(\mathcal{P}) \cap S = \emptyset$. Not to face the burden of computing yet another constant, we restrict $S$ to have size only at most $n/\log n$. (The numerical constants are, clearly, not best possible.)

**Theorem 2.16 (Connection Lemma).** For all $\alpha \in (0, 1)$ and $\delta < \delta_0$, where

$$\delta_0 = \frac{\alpha^{49}}{3^5 50^8 3000^{12}},$$

there exist two functions $r(l)$ and $\epsilon(l)$ so that for all $\mathcal{H}$, $P$ and for all integers $l$ if $(\mathcal{H}, P)$ is an $(\alpha, \delta, l, r(l), \epsilon(l))$-triad with $|V_1| = |V_2| = |V_3| = n$ sufficiently large, then there is a subgraph $P_0$ of at most $27 \sqrt{n^2/l}$ edges of $P$ such that for every ordered pair of disjoint edges $(e, f) \in (P - P_0) \times (P - P_0)$, $e \cap f = \emptyset$, and for every set $S \subset V(\mathcal{H}) \setminus (e \cup f)$ of size $|S| \leq n/\log n$, there is in $\mathcal{H} - P_0$ an $S$-avoiding hyperpath from $e$ to $f$ of length at most twelve.

**Remark 2.17.** In principle it might happen that an edge $e \in P - P_0$ is “isolated” in $\mathcal{H} - P_0$, that is, all triplets containing $e$ also contain an edge of $P_0$. The conclusion of the above theorem ensures that this is not the case. In fact, all edges $e \in P - P_0$ are mutually connected by short hyperpaths within $\mathcal{H} - P_0$. 
Figure 1. A hyperpath of length 12 from $e$ to $f$. Every 3 consecutive vertices on the path form a hyperedge.

3. Constructive reformulation

As mentioned earlier, in the case of $(d, \epsilon)$-regular graphs, it is easy to see that for every pair of vertices with at least $\epsilon n$ neighbors each, there is a short path (of length at most four) between them (see, e.g., [7] and Subsection 2.1 above). In fact, see [8], every two vertices of degree at least $16(\epsilon^2/d)n$ can be connected by a path of length at most five.

The quantification of Theorem 2.16 (note that “there exist functions $r(l)$ and $\epsilon(l)$” translates to “for all $l$ there exist $r$ and $\epsilon$”) implies the following hierarchy of constants:

$$\alpha \gg \delta, 1/l \gg 1/r, \epsilon,$$

where $\beta \gg \gamma$ means that $\gamma$ is sufficiently smaller than $\beta$, or that $\gamma$ is chosen only after $\beta$ is being fixed.

Polcyn [7], working under a comfortable assumption that $\delta \ll 1/l$, proved that most edges of $P$ can be mutually connected by hyperpaths of length at most seven. Typical edges were defined in [7] in terms of the first and second neighborhood in $\mathcal{H}$. Here, with the possibility that $\delta \gg 1/l$, to formulate a constructive version of Theorem 2.16, we need to look into the fourth neighborhood of an edge.

Let us begin by defining the first neighborhood.

**Definition 3.1.** Let $\mathcal{H}$ be a 3-uniform hypergraph and let $e = \{x, y\}$ be a pair of vertices in $V = V(\mathcal{H})$. We define the hypergraph neighborhood of $e$ to be $\Gamma_{\mathcal{H}}(e) = \{z \in V : \{z, x, y\} \in \mathcal{H}\}$. The vertices in $\Gamma_{\mathcal{H}}(e)$ will be called neighbors of $e$.

Note that in a 3-graph $\mathcal{H}$ with an underlying graph $P = P^{12} \cup P^{23} \cup P^{13}$, if $e \in P^{ij}$ then $\Gamma_{\mathcal{H}}(e) \subseteq V_k$, where $\{i, j, k\} = \{1, 2, 3\}$.

Imagine that both, $\mathcal{H}$ and $P$ are chosen at random as a result of the following 2-round experiment. First, create $P$ by tossing a coin over each pair in $(V_1 \times V_2) \cup (V_2 \times V_3) \cup (V_1 \times V_3)$ independently with the success probability $1/l$, then create $\mathcal{H}$ by selecting each triangle of $P$ with probability $\alpha$. In such a random hypergraph the expected number of triplets is $\alpha n^3/l^3$ and, for a given edge of $P$ (here we condition that $e$ has been selected), the expected value of $|\Gamma_{\mathcal{H}}(e)|$ equals $\alpha n/l^2$. It is proved in [7] that if $(\mathcal{H}, P)$ is an (deterministic) $(\alpha, \delta, l, 1, \epsilon(l))$-triad, then for almost all edges of $P$, $|\Gamma_{\mathcal{H}}(e)|$ is close to the above expectation. Here we quote without proof a minor modification of that result.
Fact 3.2 ([7]). For all $\alpha > 0$ and $\delta > 0$, there exists a function $\epsilon(l) > 0$ such that for all integers $l \geq 1$, whenever $(\mathcal{H}, P)$ is an $(\alpha, \delta, l, 1, \epsilon(l))$-triad then all but at most $7\sqrt{n^2/l}$ edges $e$ of $P^i$, $1 \leq i < j \leq 3$, satisfy the inequalities

$$n \left(\frac{1}{l} - \epsilon\right)^2 (\alpha - \delta) < |\Gamma_{\mathcal{H}}(e)| < (\alpha + \delta) \left(\frac{1}{l} + \epsilon\right)^2 n.$$ 

Definition 3.3. Let $e_1, e_2$ be edges of $P$. We say that $e_1$ reaches $e_2$ within $\mathcal{H}$ in $k$ steps and in $t$ ways if there exist at least $t$ internally disjoint hyperpaths in $\mathcal{H}$ of length $k$ from $e_1$ to $e_2$. For $t = 1$ we will skip the phrase “in $t$ ways”. For an edge $e \in P$, we denote by $\text{Four}^+(e, \mathcal{H})$ the set of those edges of $P$, which are reached from $e$ within $\mathcal{H}$ in four steps and in $\gamma_0 n$ ways, and by $\text{Four}^-(e, \mathcal{H})$ the set of all edges of $P$ which reach $e$ within $\mathcal{H}$ in four steps and in $\gamma_0 n$ ways (see Figure 2), where

$$\gamma_0 = \frac{\alpha^4}{5000l^7}.$$ 

![Figure 2. The fourth neighborhoods of $e$ ($g \in \text{Four}^-(e, \mathcal{H})$, $h \in \text{Four}^+(e, \mathcal{H})$).](image)

Let us now provide some intuition for why it is necessary to consider the fourth hypergraph neighborhood of a graph edge. Suppressing $\alpha, \delta, \epsilon$, most edges of $P$ belong to about $n/l^2$ triplets of $\mathcal{H}$ (see Fact 3.2), but any such edge $e$ can be completely cut off from the rest of $\mathcal{H}$ if no stronger assumption is made. Indeed, the total number of triplets extending triplets containing $e$ is of the order $n^2$, and clearly the removal of such a tiny fraction of triplets cannot affect the $\delta$-regularity which “controls” only sets of hyperedges of size, roughly, $n^3/l^3$.

In two steps, only about $n^2/l^4$ edges are reached from a typical edge. Most of them extend to about $n/l^2$ triplets, a total of $n^3/l^6$ – still much less than $n^3/l^3$ if $l$ is large. To estimate the number of edges reached from a typical edge in three steps, the quantity $n^3/l^6$ has to be divided to accommodate the repetitions. Among the edges $f$ reached by $e$ in three steps, only $O(n)$ share a vertex with $e$. For all other $f$, there are at most, roughly, $n/l^4$ paths from $e$ to $f$. This is because their number is bounded from above by the number of vertices forming simultaneously triangles with $e$ and $f$. Thus, there are only $n^2/l^4$ edges reached from $e$ in three steps. Again, they belong to about $n^3/l^4 \ll \delta n^3/l^3$ triplets – a quantity not under control. Hence, the shortest distance at which a typical edge can reach a substantial number of other edges is four.
Theorem 3.4 below states that, indeed, most edges have large fourth neighborhood, and, more importantly, edges with large fourth neighborhood are mutually connected by short hyperpaths.

Let us denote by \( R_0(\mathcal{H}) = R_0 \) the set of all edges of \( P \), for which

\[
\min \{ |\text{Four}^+(e, \mathcal{H})|, |\text{Four}^-(e, \mathcal{H})| \} < \frac{\alpha^4}{2000} \times \frac{n^2}{l}.
\]

**Theorem 3.4.** For all \( \alpha \in (0, 1) \) and \( \delta < \delta_0 \), where

\[
\delta_0 = \frac{\alpha^{49}}{3^{650^{5000^{12}}}},
\]

there exist two functions \( r(l) \) and \( \epsilon(l) \) such that for all \( \mathcal{H}, P \) and for all integers \( l \), if \( (\mathcal{H}, P) \) is an \( (\alpha, \delta, l, r(l), \epsilon(l)) \)-triad with \( |V_1| = |V_2| = |V_3| = n \) sufficiently large, then

(i) \( |R_0| \leq 27 \sqrt[n]{\delta n^2}/l \), and

(ii) for every ordered pair of disjoint edges \( (e, f) \in (P - R_0) \times (P - R_0) \) and for every set \( S \subset V(\mathcal{H}) \setminus (e \cup f) \) of size \( |S| \leq n/ \log n \), there is in \( \mathcal{H} \) an \( S \)-avoiding hyperpath from \( e \) to \( f \) of length at most twelve.

4. **Two lemmas and main proofs**

Theorems 2.16 and 3.4 are straightforward consequences of two technical lemmas. A subgraph \( A \) of \( P = P^{12} \cup P^{23} \cup P^{13} \) is called *framed* if for some \( 1 \leq i < j \leq 3 \), \( A \subseteq P^i \). Our first lemma needs only the assumption that \((\mathcal{H}, P)\) is an \((\alpha, \delta, l, r(l), \epsilon(l))\)-triad, where \( r(l) = 1 \) for all \( l \).

**Lemma 4.1.** For all \( c \in (0, 1) \) and \( \alpha \in (0, 1) \) and for all \( \delta < \delta_1 \), where

\[
\delta_1 = \frac{\alpha c^{12}}{3^{650^{5000^{12}}}},
\]

there exists a function \( \epsilon(l) \) so that for all \( \mathcal{H}, P \) and for all integers \( l \) if \((\mathcal{H}, P)\) is an \((\alpha, \delta, l, 1, \epsilon(l))\)-triad with \( |V_1| = |V_2| = |V_3| = n \) sufficiently large, then the following is true: For every subgraph \( P_1 \subset P \), where \( |P_1| \leq 29 \sqrt[n]{\delta n^2}/l \), and for every pair of framed subgraphs \( A \) and \( B \) of \( P - P_1 \), each of size at least \( cn^2/l \), there exist edges \( a \in A \) and \( b \in B \) and a hyperpath in \( \mathcal{H} - P_1 \) from \( a \) to \( b \) of length at most four.

Our second lemma asserts that for a typical pair \((\mathcal{H}, P)\), apart from a small set of edges \( P_0 \), all other edges of \( P \) have their fourth neighborhood substantial, even if the edges of \( P_0 \) are to be avoided. This lemma needs the whole strength of the \((\delta, r)\)-regularity.

**Lemma 4.2.** For all \( \alpha \in (0, 1) \) and \( \delta < \delta_2 \), where

\[
\delta_2 = \frac{\alpha^2}{180^2},
\]

there exist two functions \( r(l) \) and \( \epsilon(l) \) such that for all \( \mathcal{H}, P \) and for all integers \( l \) if \((\mathcal{H}, P)\) is an \((\alpha, \delta, l, r(l), \epsilon(l))\)-triad with \( |V_1| = |V_2| = |V_3| = n \) sufficiently large, then there exists \( P_0 \subset P \), \( |P_0| \leq 27 \sqrt[n]{\delta n^2}/l \), such that

\[
\min \{ |\text{Four}^+(e, \mathcal{H} - P_0)|, |\text{Four}^-(e, \mathcal{H} - P_0)| \} \geq \left( \frac{\alpha^4}{2000} \times \frac{n^2}{l} \right).
\]
for all $e \in P - P_0$.  

From Lemmas 4.1 and 4.2 we immediately derive our main result.

**Proof of Theorem 2.16.** Note that for $c = \alpha^4/3000$, $\delta_0 = \delta_1 < \delta_2$. Given $\alpha$ and $\delta < \delta_0$, let $\epsilon_1(l)$ satisfy Lemma 4.1 with $c = \alpha^4/3000$, and let functions $r(l)$, and $\epsilon_2(l)$ satisfy Lemma 4.2. We claim that Theorem 2.16 is true with the above choice of $r(l)$ and with $\epsilon(l) = \min\{\epsilon_1(l), \epsilon_2(l)\}$. 

Indeed, consider any $\mathcal{H}$, $P$ and $l$ such that $(\mathcal{H}, P)$ is an $(\alpha, \delta, l, r(l), \epsilon(l))$-triad and apply Lemma 4.2. It follows that there exists $P_0 \subset P$, $|P_0| \leq 27 \sqrt{d\delta n^2}/l$, such that (6) holds for all $e \in P - P_0$. Fix disjoint $e, f \in P - P_0$, and a set $S \subset V(\mathcal{H}) \setminus (e \cup f)$ of size $|S| \leq n/\log n$. Define $P_S = \{e \in P : S \cap e \neq \emptyset\}$ and observe that $|P_S| = O(n^2)$, and thus, for large $n$, $|P_0 \cup P_S| \leq 29 \sqrt{d\delta n^2}/l$, and

$$\min \{\left|\text{Four}^+(e, \mathcal{H} - P_0)\right|, \left|\text{Four}^-(e, \mathcal{H} - P_0)\right|\} - |P_S| \geq \frac{\alpha^4}{3000} \frac{n^2}{l}.$$  

![Figure 3. A hyperpath from e to f. (An illustration of the proof of Theorem 2.16)](image)

Since $(\mathcal{H}, P)$ is also an $(\alpha, \delta, l, \epsilon_1(l))$-triad, we may apply Lemma 4.1 with $c = \alpha^4/3000$ to

$$A = \text{Four}^+(e, \mathcal{H} - P_0) \setminus P_S, \quad B = \text{Four}^-(f, \mathcal{H} - P_0) \setminus P_S \quad \text{and} \quad P_1 = P_0 \cup P_S,$$

obtaining edges $a \in A$ and $b \in B$, and a hyperpath $\mathcal{P}_1$ in $\mathcal{H} - (P_0 \cup P_S)$ from $a$ to $b$ of length at most four. (Note that $A$ and $B$ are framed subgraphs of $P_1$.)

Let $I = V(\mathcal{P}_1) \cup f \setminus a$. Among at least $\gamma_0 n > |I \cup S|$ (for large $n$) internally disjoint hyperpaths from $e$ to $a$ in $\mathcal{H} - P_0$ choose one which is disjoint from $I \cup S$, obtaining an $S$-avoiding hyperpath $\mathcal{P}_2$ in $\mathcal{H} - P_0$ from $e$ to $b$ of length at most eight. Finally, set $J = V(\mathcal{P}_2) \setminus b$ and choose a hyperpath $\mathcal{P}_3$ in $\mathcal{H} - P_0$ from $b$ to $f$ which avoids the vertices $a$ and $b$. 

$$4 + 4 + 4 = 12$$
of \( J \cup S \). This way we obtain an \( S \)-avoiding hyperpath in \( \mathcal{H} - P_0 \) from \( e \) to \( f \) of length at most twelve (see Figure 3).

\[ \square \]

**Proof of Theorem 3.4.** Since \( R_0 \subseteq P_0 \), where \( P_0 \) is as in Lemma 4.2, part (i) follows from the estimate on \( |P_0| \). The proof of part (ii), is very similar to that of Theorem 2.16. We define \( P_S \) as before and apply Lemma 4.1 with \( c = \alpha^4/3000 \) to

\[ A = \text{Four}^+(e, \mathcal{H}) \setminus P_S, \quad B = \text{Four}^-(f, \mathcal{H}) \setminus P_S \quad \text{and} \quad P_1 = P_S, \]

obtaining edges \( a \in A \) and \( b \in B \), and a hyperpath \( \mathcal{P}_1 \) in \( \mathcal{H} - P_S \) from \( a \) to \( b \) of length at most four. Finally, we extend \( \mathcal{P}_1 \) to an \( S \)-avoiding hyperpath in \( \mathcal{H} \).

\[ \square \]

**Remark 4.3.** It will follow from the proof of Lemma 4.1 that, in fact, depending on the position of the sets \( A \) and \( B \), the promised hyperpath is precisely of length two, three or four. Consequently, depending on the position of \( e \) and \( f \), the length of a hyperpath from \( e \) to \( f \), guaranteed by Theorems 2.16 and 3.4, is precisely ten, eleven or twelve.

## 5. Short paths between large sets of edges

In this section we prove Lemma 4.1. We begin with formulating a claim from which the lemma will follow quite easily. Let \( E \) be any framed subgraph of \( P \). Further, let \( \text{First}^+(E, \mathcal{H}) \) and \( \text{Second}^+(E, \mathcal{H}) \) denote the sets of all edges \( h \in P \) reached in \( \mathcal{H} \) by an edge \( g \in E \) in one and, respectively, in two steps. Sets \( \text{First}^-(E, \mathcal{H}) \) and \( \text{Second}^-(E, \mathcal{H}) \) are defined similarly, by replacing the phrase “reached in \( \mathcal{H} \) by an edge \( g \in E \)” by “reaching in \( \mathcal{H} \) an edge \( g \in E \)". Throughout, \( ijk \) always stands for any one of the functions: 123 or 231 or 312, that is, functions which follow the cyclic ordering 1231.

**Claim 5.1.** For all \( c \in (0, 1) \) and \( \alpha \in (0, 1) \), all \( 0 < \delta < \min\{\alpha, c^{6}/50^{6}\} \) and functions \( 0 < \epsilon(l) \leq \sqrt{\delta}/(10l^{3}) \), and all integers \( l \geq 1 \), if \( (\mathcal{H}, P) \) is an \( (\alpha, \delta, l, 1, \epsilon(l)) \)-triad with \( |V_1| = |V_2| = |V_3| = n \) sufficiently large, then for all \( P_1 \subset P \) of size \( |P_1| \leq 29 \sqrt{\delta}n^{2}/l \) and for all sets \( E \subseteq P^{ij} - P_1 \) of size \( |E| \geq cn^{2}/l \),

\[
\min \left\{ |\text{First}^+(E, \mathcal{H} - P_1)|, |\text{First}^-(E, \mathcal{H} - P_1)| \right\} \geq \frac{\epsilon(n)}{6l},
\]

\[
\min \left\{ |\text{Second}^+(E, \mathcal{H} - P_1)|, |\text{Second}^-(E, \mathcal{H} - P_1)| \right\} \geq \left(1 - \frac{4\delta^{1/8}}{\sqrt{c}}\right) \frac{n^{2}}{l}.
\]

In order to derive Lemma 4.1 from Claim 5.1 we need one more simple fact about vertex-disjoint subgraphs of bipartite graphs.

**Fact 5.2.** Let \( A \) and \( B \) be two bipartite graphs with the same bipartition \( V_1 \cup V_2 \), \( |V_1| = |V_2| = n \). Then there exist \( A' \subseteq A \) and \( B' \subseteq B \) such that \( |A'| \geq (1/2)|A| - (1/2)\Delta_2(A), \) \( |B'| \geq (1/2)|B| \) and \( V(A') \cap V(B') \cap V_2 = \emptyset \), where \( \Delta_2(A) \) is the maximum degree in \( A \) among the vertices of \( V_2 \).
Proof. Let us put vertices of the set $V_2$ in two linear orders: $L_A$, ordered by their degrees in $A$ in the descending manner (ties resolved arbitrarily), and $L_B$ — the same with respect to $B$. Now include the first vertex of $L_B$ to a set $V_B$ and remove it from both orders. We repeat this step for $L_A$ and then again for $L_B$ and so on until all vertices are placed in one of the sets $V_A$ or $V_B$. (Note that $|V_A| = [n/2]$ and $|V_B| = [n/2].$)

Let us define $A'$ as the subgraph $A[V_1 \cup V_A]$ of $A$ induced by the subset of vertices $V_1 \cup V_A$, and, similarly, $B' = B[V_1 \cup V_B]$. It remains to compare the sizes of $A'$ against $A$ and $B'$ against $B$. For the latter, let us match each vertex included into $V_B$ with the one included into $V_A$ in the very next step (if $n$ is odd, the vertex included into $V_B$ last remains unmatched). Because we have started with a vertex of $V_2$ with the largest degree in $B$, its match has a smaller or equal degree in $B$, and this is true for each matched pair. Therefore, we have $|B'| \geq (1/2)|B|$. To prove that $|A'| \geq (1/2)|A| - (1/2)\Delta_2(A)$ we apply the same reasoning to the subgraph of $A$ obtained by removing all edges incident in $A$ to the first vertex of $L_B$. 

Proof of Lemma 4.1. Given $c$ and $\alpha$, let

$$\delta < \frac{\alpha c^{12}}{3^6 50^8} < \frac{\alpha (c/3)^6}{50^8}$$

and

$$\epsilon(l) = \frac{\sqrt{\delta}}{10^3 l}.$$ 

Note that $\delta < \alpha$, and

$$1 - \frac{4\delta^{1/8}}{\sqrt{c/3}} \frac{c}{12} > 1 + \epsilon(l)l,$$

the latter by inequalities $\delta^{1/8} < c \sqrt{c} / (50 \sqrt{3})$ and $\epsilon(l)l < c/50^2$.

Let $\mathcal{H}, P, l, A, B$, and $P_1$ be as in Lemma 4.1. Without loss of generality we assume that $A \subseteq P^{12}$ and will consider all three cases for $B$.

If $B \subseteq P^{13}$, apply Claim 5.1 with $E = A$ to obtain a set $A^{13} = \text{Second}^+(A, \mathcal{H} - P_1) \subseteq P^{13} - P_1$ of at least

$$\left(1 - \frac{4\delta^{1/8}}{\sqrt{c/3}}\right) \frac{n^2}{l}$$

edges. By (5) and (10), we conclude that $B \cap A^{13} \neq \emptyset$, implying the existence of a hyperpath within $\mathcal{H} - P_1$ from an edge $a \in A$ to an edge $b \in B$ of length two.

If $B \subseteq P^{12}$, we use Fact 5.2 to obtain two subgraphs $A' \subseteq A$ and $B' \subseteq B$ such that $|A'| \geq (c/3)n^2/l$ (for $n$ sufficiently large), $|B'| \geq (c/2)n^2/l$ and $V(A') \cap V(B') \cap V_2 = \emptyset$. Then by Claim 5.1 applied with $c$ replaced by $c/3$, the set $A^{13} = \text{Second}^+(A', \mathcal{H} - P_1) \subseteq P^{13} - P_1$ has cardinality at least

$$\left(1 - \frac{4\delta^{1/8}}{\sqrt{c/3}}\right) \frac{n^2}{l},$$

and taking $B^{13} = \text{First}^-(B', \mathcal{H} - P_1) \subseteq P^{13} - P_1$, by Claim 5.1 applied with $c$ replaced by $c/2$, we have

$$|B^{13}| \geq \frac{c}{12} \frac{n^2}{l}.$$
Again, by (5) and (10), we conclude that $B^{13} \cap A^{13} \neq \emptyset$. Let $zu \in B^{13} \cap A^{13}$ and let $xyzu$ and $zuw$ be hyperpaths, respectively, from $a = xy$ to $u$ and from $zu$ to $b = vu$. Note that by the disjoint choice of $A'$ and $B'$ we have $y \neq v$, and so $xyzu$ is a hyperpath within $\mathcal{H} - P_1$ from $a \in A$ to $b \in B$ of length three.

The last case is when $B \subseteq P^{2^3}$. Here also we apply Fact 5.2 to obtain two subgraphs $A' \subseteq A$ and $B' \subseteq B$ such that $|A'| \geq (c/3)n^2/l$, $|B'| \geq (c/2)n^2/l$ and $V(A') \cap V(B') \cap V_2 = \emptyset$. (Technically, we identify for a moment sets $V_1$ and $V_3$ to treat $A$ and $B$ as two bipartite graphs on the same vertex set.) By Claim 5.1 applied with $c$ replaced by $c/3$, the set $A^{13} = \text{Second}^+(A', \mathcal{H} - P_1) \subseteq P^{13} - P_1$ consists of at least
\[
\left(1 - \frac{4\delta^{1/8}}{\sqrt{c/3}}\right) \frac{n^2}{l}
\]
edges. Hence, taking $B^{13} = \text{Second}^-(B', \mathcal{H} - P_1) \subseteq P^{13} - P_1$, by Claim 5.1 applied with $c$ replaced by $c/2$, we get
\[
|B^{13}| \geq \left(1 - \frac{4\delta^{1/8}}{\sqrt{c/2}}\right) \frac{n^2}{l}.
\]
Again, by (5), (10) and (9), we conclude that $B^{13} \cap A^{13} \neq \emptyset$. Let $zu \in B^{13} \cap A^{13}$ and let $xyzu$ and $zuw$ be hyperpaths, respectively, from $a = xy$ to $u$ and from $zu$ to $b = vu$. Note that by the disjoint choice of $A'$ and $B'$ we have $y \neq v$, and so $xyzuw$ is a hyperpath within $\mathcal{H} - P_1$ from $a \in A$ to $b \in B$ of length four (see Figure 4).

![Figure 4. An illustration of the last case of the proof of Lemma 4.1.](image)

It remains to prove Claim 5.1. We first show a simple but crucial fact which will be applied twice in the proof of Claim 5.1.

**Fact 5.3.** For any real $\alpha, \delta \in (0, 1)$, integer $l$ and $\epsilon \leq \sqrt{\delta}/(10l^3)$, let $(\mathcal{H}, P)$ be an $(\alpha, \delta, l, 1, \epsilon)$-triad. Let, further, $A \subseteq V_j$, $B \subseteq V_i$ and $Q \subseteq P$ be such that $|P^{h_{\mathcal{H}} - Q}| \leq 29 \sqrt{\delta}n^2/l$, $|A| \geq an$, $|B| \geq bn$, and every vertex of $A$ has in $Q$ at least $\beta|B|/l$ neighbors in $B$ and at least $\gamma n/l$ neighbors in $V_k$. If $\min\{\gamma, b\beta\} > \epsilon l$ and $ab\beta \delta \geq 43 \sqrt{\delta}$ then $|\mathcal{H} \cap \text{Tr}(Q)| > 2(\alpha - \delta) \delta n^3/l^3$. 
Proof. By the $(1/l, \epsilon)$-regularity of $P^k$, we have
\[
|\text{Tr}(Q \cup P^k)| \geq \sum_{y \in A} \deg_Q(y, B) \deg_Q(y, V_i) \left( \frac{1}{l} - \epsilon \right) \geq \frac{9}{10} ab \beta y \frac{n^3}{l^3}.
\]
On the other hand, setting $Q' = P^j \cup P^k \cup (P^k - Q)$, by Corollary 2.4, we have
\[
|\text{Tr}(Q')| < 29 \sqrt{\delta} (1.21) \frac{n^3}{l^3} + 4 \epsilon n^3 < 36 \sqrt{\delta} \frac{n^3}{l^3},
\]
and so, by our assumptions and Fact 2.6, we may estimate
\[
|\text{Tr}(Q)| \geq |\text{Tr}(Q \cup P^k)| - |\text{Tr}(Q')| \geq \left( \frac{9}{10} - 43 \sqrt{\delta} - 36 \sqrt{\delta} \right) \frac{n^3}{l^3} > 2 \delta \frac{n^3}{l^3} > \delta |\text{Tr}(P)|.
\]
Therefore, by the $(\alpha, \delta, 1)$-regularity of $H$,
\[
d_H(Q) = \frac{|H \cap \text{Tr}(Q)|}{|\text{Tr}(Q)|} > \alpha - \delta. \quad \square
\]

Proof of Claim 5.1. By symmetry, it is enough to prove only that $|\text{First}^+(E, H - P_1)| \geq (c/6)n^2/l$ and similarly, that $|\text{Second}^+(E, H - P_1)| \geq \left(1 - 4\delta^{1/8}/\sqrt{c} \right)n^2/l$. Let us fix $\alpha$ and $c, 0 < \alpha, c < 1$, and let
\[
\delta < \min \left\{ \alpha, \frac{c\delta}{50^8} \right\}.
\]
Further, with $\delta$ given above, let for all $l$
\[
\epsilon(l) \leq \frac{\sqrt{\delta}}{10l^5} < \frac{\delta^{1/4}}{l\sqrt{c}} < \frac{c}{120l}.
\]
Let $H, P$, and $l$ be as in Claim 5.1. Set $\epsilon = \epsilon(l)$ for convenience and fix $1 \leq i < j \leq 3$. Let $E$ be a set of at least $cn^2/l$ edges of $P^j - P_1$. Define
\[
E_1 = \{yz \in P^j : \text{xyz} \in H \text{ and } xy \in E \text{ and } xz \notin P_1 \text{ for some } x \in V_i \}
\]
and assign to each edge $yz$ of $E_1$ one (arbitrary) vertex $x = x_{yz} \in V_i$ which together with $yz$ satisfies the conditions in the definition of $E_1$. Finally, let
\[
E_2 = \bigcup_{yz \in E_1 - P_1} \{zw \in P^k : \text{w} \neq x_{yz} \text{ and } yzw \in H \text{ and } yw \notin P_1 \}.
\]
Note that $E_1 - P_1 = \text{First}^+(E, H - P_1)$, and that by avoiding $w = x_{yz}, E_2 - P_1 \subseteq \text{Second}^+(E, H - P_1)$.

Observation 5.4. Trivially, if $xy \in E$, $yz \in P^j - E_1$ and $xz \notin P_1$ then $xyz \notin H$. Similarly, but more subtly, if $yz \in E_1 - P_1, zw \in P^k - E_2$, and $yw \notin P_1$, then $yzw \notin H$ unless $w = x_{yz}$, which implies that the edges of $E_1 - P_1, P^k - E_2$ and $P^j - P_1$ span at most $|E_1 - P_1|$ hyperedges in $H$. 
Using these observations and Fact 5.3 we will first show that a significant fraction of vertices $y \in V_j$ have large (close to $n/l$) neighborhood in $E_1$, and so subgraph $E_1 - P_1$ is large. Then we will argue that most vertices of $V_k$ have large degree in $E_2$, meaning that the set $E_2$ must be very large (close to $n^2/l$), and so must be $E_2 - P_1$.

Let

$$L_0 = \left\{ y \in V_j : \deg_{P_{jk}}(y) < \left( \frac{1}{l} - \epsilon \right)n \right\}.$$  

By (2) with $A = V_k$, we have $|L_0| \leq \epsilon n$. Next, let us consider the set

$$L = \left\{ y \in V_j - L_0 : \deg_{E_1}(y) \geq \frac{cn}{2l} \right\}.$$

Observe that $|L| \geq \frac{cn}{3}$. Indeed, otherwise, using (1) and (12), we obtain a contradiction

$$|E| < |L| \left( \frac{1}{l} + \epsilon \right) + |L_0| \left( \frac{1}{l} - \epsilon \right) + \epsilon n^2 + n \frac{cn}{2l} < \frac{cn^2}{l}.$$

We proceed with the following fact. Set $\overline{E}_1 = P_{jk} - E_1$ and

$$L' = \left\{ y \in L : \deg_{E_1}(y) > \frac{7\delta^{1/4}n}{\sqrt{c} l} \right\}.$$

**Fact 5.5.**

$$|L'| \leq 13 \frac{\delta^{1/4}n}{\sqrt{c}}$$

**Proof.** Assume $|L'| > 13(\delta^{1/4}/\sqrt{c})n$ and apply Fact 5.3 with $A = L'$, $B = V_i$ and so $b = 1$, $\beta = c/2$, $a = 13\delta^{1/4}/\sqrt{c}$, and $\gamma = 7\delta^{1/4}/\sqrt{c}$ to the 3-partite subgraph $Q = Q^{ij} \cup Q^{jk} \cup Q^{ki}$, where

$$Q^{ij} = E[V_i, L'], \\
Q^{jk} = \overline{E}_1[L', V_k], \\
Q^{ki} = P_{ki} - P_1.$$

As $\min\{\gamma, b\beta\} > el$ and $ab\beta\gamma > 43 \sqrt{d}$, it follows that $\mathcal{H} \cap \text{Tr}(Q) \neq \emptyset$. However, by the construction of $Q$ (see Observation 5.4) we have $\mathcal{H} \cap \text{Tr}(Q) = \emptyset$. This contradiction ends the proof of Fact 5.5. \qed

We set $L'' = L - L'$. By (11),

$$|L''| = |L| - |L'| \geq \frac{1}{3}cn - 13\frac{\delta^{1/4}n}{\sqrt{c}} > \frac{1}{4}cn > \epsilon n.$$  

Note that every vertex $y \in L''$ has

$$\deg_{E_1}(y) > \frac{n}{l} - \epsilon n - 7\frac{\delta^{1/4}n}{\sqrt{c} l},$$

and thus, by (11), (12) and (13),

$$|E_1| \geq |L''|(\frac{n}{l} - \epsilon n - 7\frac{\delta^{1/4}n}{\sqrt{c} l}) > \frac{c n^2}{5 l}.$$
To complete the proof of the inequality (7) we count the number of edges in $E_1 - P_1$

$$|E_1 - P_1| > \frac{c n^2}{5l} - 29 \sqrt{\frac{\delta n^2}{l}} > \frac{c n^2}{6l},$$

the latter by (11).

We continue with the proof of the inequality (8). Recall our notation $\deg_{G}\(v\)$ and $\deg_{G}(v,U)$ defined in Subsection 2.1. Let $\bar{E}_2 = P^{ki} - E_2$. Note that by the $(1/l, \epsilon)$-regularity of $P^{ji}$ and $P^{ki}$, Fact 2.3 and (13), the set

$$L_0 = \left\{z \in V_k : \deg_{E_2}(z, L'') < \left(\frac{1}{l} - \epsilon\right)|L''| \text{ or } \deg_{E_2}(z) < \left(\frac{1}{l} - \epsilon\right)n\right\}$$

has size $|L_0| \leq 2en$. Next, let us consider the set

$$L_1 = \left\{z \in V_k : \deg_{E_1}(z, L'') > 7\frac{\delta^{1/8}}{l}|L''|\right\}.$$

Since each vertex of $L''$ has in $E_1$ degree at most $7(\delta^{1/4}/\sqrt{c})(n/l)$, a simple, double counting argument shows that $|L_1| \leq (\delta^{1/4}/\sqrt{c})n$. Further, let

$$L_2 = \left\{z \in V_k : \deg_{P_i}(z, L'') > 116\frac{\delta^{1/4}n}{\sqrt{cl}}|L''|\right\}.$$

Clearly, $|L_2| < (\delta^{1/4}/\sqrt{c})n$, since otherwise $|P_1| > 29 \sqrt{\delta n^2}/l$ – a contradiction. Set $L = V_k \setminus (L_0 \cup L_1 \cup L_2)$ and define

$$L' = \left\{z \in L : \deg_{E_2}(z) > 9\frac{\delta^{1/4} n}{\sqrt{c l}}\right\}.$$

Observe, by (11) and (12), that for all $z \in L$ (and thus for all $z \in L'$) we have

$$\deg_{E_1-P_1}(z, L'') > \left(\frac{1}{l} - \epsilon - 7\frac{\delta^{1/8}}{l} - 116\frac{\delta^{1/4} n}{\sqrt{cl}}\right)|L''| > \frac{4}{5l}|L''|.$$

**Fact 5.6.**

$$|L'| \leq 24\frac{\delta^{1/4} n}{\sqrt{c}}$$

**Proof.** The proof of this fact is very similar to the proof of Fact 5.5. We will argue that the inequality $|L'| > 24(\delta^{1/4}/\sqrt{c})n$ contradicts a conclusion of Fact 5.3. Define a 3-partite subgraph $Q = Q^{ij} \cup Q^{ik} \cup Q^{ki}$ as follows:

$$Q^{ij} = P^{ij} - P_1,$$

$$Q^{ik} = E_1[L', L''] - P_1,$$

$$Q^{ki} = \bar{E}_2[L', V_i].$$

By the construction of $Q$ and Observation 5.4 we have $|\mathcal{H} \cap \text{Tr}(Q)| \leq |E_1| = O(n^2)$. Apply Fact 5.3 with $A = L'$, $B = L''$ (and so $b = e/4$), $\beta = 4/5$, $a = 24\delta^{1/4}/\sqrt{c}$, and $\gamma = 9\delta^{1/4}/\sqrt{c}$ to yield $|\mathcal{H} \cap \text{Tr}(Q)| = \Omega(n^2)$. For large enough $n$, this is a contradiction which ends the proof of Fact 5.6. □
To complete the proof of the inequality (8), set
\[ \mathcal{L}'' = \mathcal{L} \setminus \mathcal{L}' = V_2 \setminus (\mathcal{L}_0 \cup \mathcal{L}_1 \cup \mathcal{L}_2 \cup \mathcal{L}') \]
and note that for every vertex \( z \in \mathcal{L}'' \), by (12), we have
\[
\deg_{E_2}(z) \geq \frac{n}{l} - e n - 9 \frac{\delta^{1/4}}{\sqrt{c}} \frac{n}{l} > (1 - 10 \frac{\delta^{1/4}}{\sqrt{c}}) \frac{n}{l}.
\]
Note also that all the exceptional sets \( \mathcal{L}', \mathcal{L}_0, \mathcal{L}_1 \) and \( \mathcal{L}_2 \) contain together less than \( 2(\delta^{1/8}/\sqrt{c}) n \) vertices and therefore \( |\mathcal{L}''| > (1 - 2\delta^{1/8}/\sqrt{c}) n \). Thus, by (11),
\[
|E_2| \geq |\mathcal{L}''| \left( 1 - 10 \frac{\delta^{1/4}}{\sqrt{c}} \right) \frac{n}{l} > \left( 1 - 3 \frac{\delta^{1/8}}{\sqrt{c}} \right) \frac{n^2}{l},
\]
hence
\[
|E_2 - P| > \left( 1 - 3 \frac{\delta^{1/8}}{\sqrt{c}} \right) \frac{n^2}{l} - 2.9 \sqrt{\delta n} \frac{n^2}{l} > \left( 1 - 4 \frac{\delta^{1/8}}{\sqrt{c}} \right) \frac{n^2}{l}.
\]
\[ \square \]

6. THE FOURTH NEIGHBORHOOD

In this section we prove Lemma 4.2. Let us begin with some heuristic. We call an edge \( e \in \mathcal{H} \) \( \mathcal{H} \)-good, or just good if, say, \( |\Gamma_{\mathcal{H}}(e)| \geq (2/9)an/|l|^2 \). We call an edge bad if it is not good. As proved in [7] (see Fact 3.2 above), for most edges \( e \) of \( P \) we have \( |\Gamma_{\mathcal{H}}(e)| \sim an/|l|^2 \), so most edges are good, but unfortunately, some of these good edges may have small fourth, and even second neighborhood. Indeed, it might happen that for a good edge \( e = xy \), whenever \( xyz \in \mathcal{H} \) then \( yz \) has a very small neighborhood.

To find a large subset of good edges \( e \) with large fourth neighborhoods \( \text{Four}^+(e, \mathcal{H}) \) and \( \text{Four}^-(e, \mathcal{H}) \), one could argue as follows. Suppose that the set of bad edges has size \( \rho n^2 \). Then, for each \( i = 1, 2, 3 \), at most \( \sqrt{\rho} n \) vertices of \( V_i \) are incident to at least \( \sqrt{\rho} n \) bad edges (let us call these vertices bad), and, provided \( \sqrt{\rho} \ll 1/|l|^2 \), one could start at a good edge with good endpoints and move four steps, avoiding both, bad edges and bad vertices. The problem is that Fact 3.2 yields only \( \rho \) of order \( 1/|l| \) – too large for our needs.

To get around this problem we will find a sub-hypergraph \( \mathcal{H}' \subseteq \mathcal{H} \) with much less bad edges. This sounds paradoxical, since removing hyperedges can only decrease \( |\Gamma_{\mathcal{H}}(e)| \). Note, however, that edges \( e \) with \( \Gamma_{\mathcal{H}}(e) = \emptyset \) are not so bad – there is no way to get to them! Let us call them \( \mathcal{H} \)-dead. To distinguish between \( \mathcal{H} \)-dead and other bad edges, we will alter our previous definition and call an edge \( e \in P \) \( \mathcal{H} \)-bad if
\[
0 < |\Gamma_{\mathcal{H}}(e)| < \frac{2}{9} \frac{n}{|l|^2}.
\]
So, for any \( \mathcal{H}' \subseteq \mathcal{H} \), every edge of \( P \) is either \( \mathcal{H}' \)-good or \( \mathcal{H}' \)-bad or \( \mathcal{H}' \)-dead. Let us denote these three subgraphs by \( G_{\mathcal{H}'}, B_{\mathcal{H}'} \) and \( D_{\mathcal{H}'} \). For technical reasons we distinguish also a class \( F_0 \) of atypical edges of \( P \). For all \( 1 \leq i < j \leq 3 \), an edge \( e \in P^i \) belongs to the subgraph \( F_0 \), if either it is not typical or at least one of its ends is not typical in \( P^j \) (see Definition 2.2). Note, that by Fact 2.3 and Corollary 2.4, \( |F_0| \leq 24\epsilon n^2 \).
Claim 6.1. For all $\alpha \in (0, 1)$ and $\delta < \alpha/\sqrt{2}$ there exist two functions $r(l)$ and $\epsilon(l)$ so that for all $\mathcal{H}$, $P$ and for all integers $l$ if $(\mathcal{H}, P)$ is an $(\alpha, \delta, l, r(l), \epsilon(l))$-triad then there exists a sub-hypergraph $\mathcal{H}' \subseteq \mathcal{H}$ such that $|B_{\mathcal{H}}| \leq \delta n^2/l^4$, $|D_{\mathcal{H}}| \leq 22 \sqrt{\delta} n^2/l$ and $F_0 \subseteq D_{\mathcal{H}}$.

Proof of Lemma 4.2. With given $\alpha$ and

$$\delta < \delta_2 = \frac{\alpha^2}{180^2} < \frac{\alpha}{92},$$

let $r(l)$ and $\epsilon_1(l)$ be such that Claim 6.1 holds. Set

$$\epsilon(l) = \min\left\{\epsilon_1(l), \frac{\alpha^4}{20,000}\right\}.$$

We will prove Lemma 4.2 with this choice of functions $r(l)$ and $\epsilon(l)$. Given integer $l$, let a pair $(\mathcal{H}, P)$ be an $(\alpha, \delta, l, r, \epsilon)$-triad, where $r = r(l)$ and $\epsilon = \epsilon(l)$. Further, let $\mathcal{H}$ be as in Claim 6.1, let

$$V^* = \left\{v \in V : \deg_{B_{\mathcal{H}}}(v) \geq \sqrt{\delta} \frac{n}{l^2}\right\},$$

$$G^*_{\mathcal{H}} = \left\{e \in G_{\mathcal{H}} : e \cap V^* \neq \emptyset\right\},$$

and let $P_0 = B_{\mathcal{H}} \cup D_{\mathcal{H}} \cup G^*_{\mathcal{H}}$.

Note that $F_0 \subseteq P_0$. It remains to prove two facts about $P_0$.

Fact 6.2.

$$|P_0| \leq 27 \sqrt{\delta} \frac{n^2}{l}.$$

Proof. To prove this fact, note that $|V^*| \leq 2 \sqrt{\delta} n/l^2$, and so $|G^*_{\mathcal{H}}| \leq 2n|V^*| \leq 4 \sqrt{\delta} n^2/l^2$. Therefore

$$|P_0| \leq |B_{\mathcal{H}}| + |D_{\mathcal{H}}| + |G^*_{\mathcal{H}}| \leq \delta \frac{n^2}{l^2} + 22 \frac{\sqrt{\delta} n^2}{l} + 4 \sqrt{\delta} \frac{n^2}{l} \leq 27 \sqrt{\delta} \frac{n^2}{l}.$$

Fact 6.3. For every edge $e \in P - P_0$ the inequality (6) holds.

Proof. By symmetry, we will only prove that $|\text{Four}^+(e, \mathcal{H} - P_0)| \geq \left(\alpha^4/2000\right)n^2/l$. Without loss of generality we may assume that $e = xy \in P^{12} - P_0$, where $x \in V_1$. Then, by our choice of $\delta$, the set of vertices $z$, such that $xyz \in \mathcal{H}$ and $yz, xz \notin P_0$, has size at least

$$\left(\frac{2}{9}\right)\alpha \frac{n}{l^2} - 4 \sqrt{\delta} \frac{n}{l^2} > \frac{\alpha n}{5l^2},$$

where the deletion takes care of all $z \in V^*$, as well as all $z$ with $yz$ or $xz$ in $B_{\mathcal{H}}$ (clearly, $yz$ and $xz$ cannot be $\mathcal{H}$-dead). Thus, $xyz \in \mathcal{H} - P_0$. For each such $z$, the edge $yz$ belongs in turn to at least $an/(5l^2)$ triplets $yzw \in \mathcal{H}'$ with $w \in V_1 \setminus \{x\}$, $yw \in P^{12} - P_0$ and $zw \in P^{13} - P_0$. So, altogether there are at least $\alpha^2 n^2/(25l^4)$ edges of $P^{13} - P_0$ reached (within $\mathcal{H} - P_0$) in two steps by $e$. Repeating this argument again we obtain at least $\alpha^4 n^4/(625l^8)$ hyperpaths $xyzuvw \in \mathcal{H} - P_0$ of length four originating at $e = xy$.

Let us estimate, by counting repetitions, how many different edges $uv \in P^{23}, u \in V_2, v \in V_3$, are indeed reached by $e$ in four steps (and in many ways). Consider an auxiliary
bipartite graph $C = (X, Y, E_C)$, where $X = E(P^{13})$, $Y = E(P^{23})$, and $\{zw \in X, uv \in Y\} \in E_C$ if $xyzwuv$ is a hyperpath in $\mathcal{H} - P_0$. Hence, $|E_C| \geq \alpha^4 n^4/(625l^8)$.

Every hyperpath $xyzwuv$ must satisfy that
\[ z \in N_{p^{23}}(u, N_p(xy)) \text{ and } w \in N_p(uv, N_{p^{12}}(y)) \]
(see Figure 5).

\[ \Delta^* = \max_{e \in Y \setminus Y_1} \deg_{C}(e). \]

Then
\[ \Delta^* \leq \Delta_0 = \frac{20 n^2}{19 l^7}. \]

Let
\[ Y_2 = \left\{ uv \in Y : \deg_{C}(uv) \geq \frac{1}{2} \frac{|E_C|}{|P^{23}|} \right\}, \]
and $Y_3 = Y \setminus (Y_1 \cup Y_2)$. We have
\[ |E_C| \leq |Y_1| n^2 + |Y_2| \Delta_0 + |Y_3| \frac{|E_C|}{2|P^{23}|} \leq 4en^4 + |Y_2| \Delta_0 + \frac{|E_C|}{2}. \]

Therefore, by our choice of $\epsilon$,
\[ |Y_2| \geq \left( \frac{\alpha^4 n^4}{1250l^8} - 4en^4 \right) \frac{1}{\Delta_0} > \frac{\alpha^4 n^2}{2000l} \]

Another words, at least $\frac{\alpha^4 n^2}{2000l}$ edges $uv \in P^{23}$ can be reached from $e$ by no less than
\[ \frac{1}{2} \frac{|E_C|}{|P^{23}|} > \frac{\alpha^4 n^2}{2500l^7} \]
hyperpaths of length four, or equivalently, via that many edges $zw \in P^{13}$. It is easy to see
that among these edges there is a matching of size at least

$$\frac{\alpha^4 n}{5000^7}.$$ 

\[ \square \]

Proof of Claim 6.1.

Given $\alpha$, let

$$\delta < \frac{\alpha}{9^2},$$

and let $\epsilon_i(l)$ be such that Fact 3.2 holds with above $\alpha$ and $\delta$. Further, let for all $l$,

$$r(l) = 32l^3$$

and

$$\epsilon(l) = \min\left\{\epsilon_i(l), \frac{\delta}{24\beta}\right\}.\tag{15}$$

We will prove Claim 6.1 with this choice of $\delta$, $r(l)$ and $\epsilon(l)$. Given integer $l$, let a pair $(\mathcal{H}, P)$ be an $(\alpha, \delta, l, r, \epsilon)$-triad, where $r = r(l)$ and $\epsilon = \epsilon(l)$.

We will define a process of deleting hyperedges which after finitely many rounds will arrive at a sub-hypergraph $\mathcal{H}'$ of $\mathcal{H}$ satisfying the conclusions of Claim 6.1. Recall that for an arbitrary hypergraph $\mathcal{H}$ and a graph $G$, we denote by $\mathcal{H} - G$ the sub-hypergraph of $\mathcal{H}$ obtained by removing all hyperedges containing at least one edge of $G$.

The initial step of the procedure isolates all edges of $F_0$. Set $\mathcal{H}_1 = \mathcal{H} - F_0$. Clearly, for each $e \in F_0$, we have $\Gamma_{\mathcal{H}_1}(e) = \emptyset$ and so $e$ is $\mathcal{H}_1$-dead.

In each next round we similarly “kill” edges of $P$ which are bad in the current sub-hypergraph. For technical reasons these rounds take cyclically care of the edges of $P^{12}$, $P^{23}$, and $P^{13}$. For each $s = 1, 4, 7, \ldots$, let

$$F_s = \{e \in P^{12} : e \text{ is } \mathcal{H}_s\text{-bad}\}, \quad \mathcal{H}_{s+1} = \mathcal{H}_s - F_s,$$

$$F_{s+1} = \{e \in P^{23} : e \text{ is } \mathcal{H}_{s+1}\text{-bad}\}, \quad \mathcal{H}_{s+2} = \mathcal{H}_{s+1} - F_{s+1},$$

$$F_{s+2} = \{e \in P^{13} : e \text{ is } \mathcal{H}_{s+2}\text{-bad}\}, \quad \mathcal{H}_{s+3} = \mathcal{H}_{s+2} - F_{s+2}.$$

In each operation of the type $\mathcal{H}_{s+1} = \mathcal{H}_s - F_s$ we remove all hyperedges which contain $\mathcal{H}_s$-bad edges $e$ of $P^{12}$, $P^{23}$ or $P^{13}$. Thus, those edges become $\mathcal{H}_{s+1}$-dead and therefore will never become bad again. It follows that all sets $F_s$ are disjoint, and, in particular, for $s \geq 1$, $F_s \cap F_0 = \emptyset$.

Our immediate goal is to estimate $\sum_{s=1}^{r} |F_s|$. Let us define 3-partite subgraphs $Q_s$ of $P$, $s = 1, 2, \ldots, r$, as follows: If $F_s \subset P^j$ then

$$Q_s = F_s \cup \left( P^{ik} - \bigcup_{t=1}^{s} F_t \right) \cup \left( P^{jk} - \bigcup_{t=1}^{s} F_t \right),$$

Set $Q = (Q_1, \ldots, Q_r)$. 


Observe that for slightly enlarged subgraphs $\overline{Q}_s = F_s \cup P^{ik} \cup P^{jk}$ (where $F_s \subset P^{ij}$), we have, by (15) and the fact that $F_s \cap F_0 = \emptyset$,
\[
|\text{Tr}(\overline{Q}_s)| \geq |F_s|n \left( \frac{1}{l} - \epsilon \right)^2 \geq \frac{3}{4} n^2 |F_s|.
\]
Trivially,
\[
\bigcup_{s=1}^r \text{Tr}(Q_s) \subseteq \bigcup_{s=1}^r \text{Tr}(\overline{Q}_s),
\]
but the reverse inclusion is also true. Indeed, for $xyz \in \text{Tr}(\overline{Q}_s)$ set $t_0 = \min\{t : \{xy, xz, yz\} \cap F_t \neq \emptyset\}$. Then $t_0 \leq s$ and $xyz \in \text{Tr}(Q_{t_0})$. Moreover, because the sets $F_s$ are disjoint, we have
\[
\left| \bigcup_{s=1}^r \text{Tr}(\overline{Q}_s) \right| \geq \frac{1}{3} \sum_{s=1}^r |\text{Tr}(\overline{Q}_s)|.
\]
Hence,
\[
\left| \bigcup_{s=1}^r \text{Tr}(Q_s) \right| = \left| \bigcup_{s=1}^r \text{Tr}(\overline{Q}_s) \right| \geq \frac{1}{4} \sum_{s=1}^r |F_s|.
\]
On the other hand, however, by the definition of an $\mathcal{H}_s$-bad edge, for all $s \leq r$,
\[
|\mathcal{H} \cap \text{Tr}(Q_s)| < |F_s| \frac{2}{9} \alpha \frac{n}{l^2},
\]
forcing
\[
d_{\mathcal{H}}(Q) < \frac{8}{9} \alpha,
\]
where $d_{\mathcal{H}}(Q)$ is defined in (3). Therefore, by the $(\alpha, \delta, r)$-regularity of $\mathcal{H}$,
\[
\left| \bigcup_{s=1}^r \text{Tr}(Q_s) \right| \leq \delta |\text{Tr}(P)|,
\]
since otherwise $d_{\mathcal{H}}(Q) > \alpha - \delta \geq (8/9)\alpha$. This inequality together with (15), (16) and Fact 2.6 implies that
\[
\sum_{s=1}^r |F_s| \leq 4 \left| \bigcup_{s=1}^r \text{Tr}(Q_s) \right| \frac{\epsilon^2}{n} < 8 \delta \frac{n^2}{l^2}.
\]
Thus, more than a half of the sets $F_s$, $s \leq r$, have size $|F_s| \leq 16 \delta n^2 / l^2$, and so two consecutive sets must be such, that is, there exists an index $s \leq r - 2$, such that
\[
\max \{|F_{s+1}|, |F_{s+2}|\} \leq 16 \delta \frac{n^2}{l^2} = \frac{1}{2} \delta \frac{n^2}{l^2}.
\]
Let $s_0$ be the smallest index $s$ with this property.

Without loss of generality we may assume that $F_{s_0} \subset P^{12}$. Set $\mathcal{H}' = \mathcal{H}_{s_0+1}$. Observe that there is no $\mathcal{H}'$-bad edge in the graph $P^{12}$, while in each $P^{23}$ and $P^{13}$ we have at most $(1/2)\delta n^2 / l^2$ $\mathcal{H}'$-bad edges. In fact, the set of $\mathcal{H}'$-bad edges is the union of $F_{s_0+1}$ ($\mathcal{H}'$-bad edges in $P^{23}$) and a subgraph of $F_{s_0+2}$ ($F_{s_0+2}$ may contain $\mathcal{H}_{s_0+2}$-bad edges which were not $\mathcal{H}_{s_0+1}$-bad).
As for the $H'$-dead edges, these are exactly the edges in $\bigcup_{j=0}^{s_0} F_j$ plus all the edges $e \in P$ which were originally dead, that is, which had $\Gamma_{H'}(e) = \emptyset$. We have already estimated $|\bigcup_{j=1}^{s_0} F_j|$ in (18), while, by (15), $|F_0| < 24en^2 < \sqrt{\delta n^2}/l$. Finally, by Fact 3.2, there are no more than $21 \sqrt{\delta n^2}/l$ originally dead edges. Therefore we have

$$|D_{H'}| < \left| \bigcup_{j=1}^{s_0} F_j \right| + |F_0| + |D_{H'}| < 8\delta \frac{n^2}{l} + \delta \frac{n^2}{l} + 21 \sqrt{\delta \frac{n^2}{l}} < 22 \sqrt{\delta \frac{n^2}{l}}$$

Hence, Claim 6.1 is proved. \hfill \Box

7. Applications

7.1. Long hyperpaths. The “Blow-up Lemma” of Komlós, Sárközy and Szemerédi [5] states that with a suitable choice of parameters every $s$-partite graph $G$ with $s$-partition $V_1 \cup \cdots \cup V_s$ in which all bipartite subgraphs $G[V_i, V_j]$ are $(d, \epsilon)$-regular contains all bounded degree $s$-partite graphs $G'$ with $s$-partition $V'_1 \cup \cdots \cup V'_s$, where for all $i = 1, \ldots, s$, $V'_i \subseteq V_i$, $|V'_i| < (1 - f(\epsilon))|V_i|$.

So far no analogous results exist for 3-uniform hypergraphs. As a first step toward a hypergraph “Blow-up Lemma”, we derive from Corollary 3.4 a simple consequence which establishes the existence of an almost Hamiltonian hyperpath in a quasi-random 3-graph.

**Proposition 7.1.** For all $\alpha \in (0, 1)$ and $\delta < (\delta_0/4)^4$, where $\delta_0$ is as in Theorem 3.4, there exist two functions $r(l)$ and $\epsilon(l)$ such that for all $H$, $P$ and for all integers $l$ if a pair $(H, P)$ is an $(\alpha, \delta, l, r(l), \epsilon(l))$-triad with $|V| = n$ sufficiently large, then there is in $H$ a hyperpath of length at least $(1 - \delta^{1/4})n$.

**Proof.** Given $\alpha$, let $\delta < (\delta_0/4)^4$ and let $r = r(l)$ and $\epsilon = \epsilon(l)$ be ensured by Theorem 3.4. Set $\epsilon = \epsilon(l) = \delta^{1/4} \epsilon(l)$. Observe, that

$$27 \sqrt{4\delta^{1/4}} < 27 \sqrt{\delta_0} < \frac{a^4}{2000}.$$  

Let a pair $(H, P)$ be an $(\alpha, \delta, l, r, \epsilon)$-triad. Suppose, that no hyperpath in $H$ has length $(1 - \delta^{1/4})n$. For a hyperpath $Q$, let $H'_Q$ be the sub-hypergraph of $H$, obtained by deleting from $H$ all, but the last four vertices of the path $Q$ (if $|V(Q)| < 4$, then we set $H'_Q = H$).

Let us fix an arbitrary edge $e = [x, y] \in P - R_0$ and let $Q$ be the longest hyperpath in $H$ originating at $e$ (in the cyclic order $V_1 \to V_2 \to V_3 \to V_1$) and such that its other endpair $f \in P - R_0(H'_Q)$. It follows trivially from the definition of the set $R_0(H)$ that $Q$ has at least four vertices. Let us denote the last four vertices of $Q$ by $x_{-3}, x_{-2}, x_{-1}, x_0$.

Since $|V(Q)| < (1 - \delta^{1/4})n$, the sub-hypergraph $H''_Q = H - V(Q)$ has at least $\delta^{1/4}n$ vertices. Moreover, since $Q$ traverses the sets $V_1, V_2, V_3$ in the cyclic order, the sizes of the sets $V(H''_Q) \cap V_i$, $i = 1, 2, 3$, differ from each other by at most one. Hence (see [9], Fact 4.2), the pair $(H''_Q, P[V(H''_Q)])$ is an $(\alpha, 4\delta^{1/4}, l, r, \epsilon/\delta^{1/4})$-triad. Note that $|V(H'_Q)| = |V(H''_Q)| + 4$, $4\delta^{1/4} < \delta_0$ and $\epsilon/\delta^{1/4} = \epsilon(l)$. Therefore, by Theorem 3.4,

$$|R_0(H''_Q)| \leq 27 \sqrt{4\delta^{1/4}} \frac{|V(H''_Q)|/3|^2}{l} \leq 27 \sqrt{4\delta^{1/4}} \frac{|V(H'_Q)|/3|^2}{l}.$$
On the other hand, by the definition of $R_0(H'_Q)$, we know that the edge $f = \{x_{-1}, x_0\}$ reaches in four steps at least
\begin{align*}
\frac{\alpha^4 |V(H'_Q)|/3^2}{2000} > 27 \sqrt{4 \delta^{1/4}} |V(H'_Q)|/3^2 + 2n \geq |R_0(H'_Q)| + 2n
\end{align*}
other edges of $P[|V(H'_Q)|]$ (the term $2n$ takes care of all edges with at least one endpoint in $x_{-2}$ or $x_{-3}$; the first inequality follows from (19) for large $n$). Therefore, there is at least one edge $f' = \{x_3, x_4\} \in P[|V(H''_Q)|] - R_0(H''_Q)$, reached by $f$ in $H'_Q$ by at least three (in fact, many more) internally disjoint hyperpaths of length four of the form $x_{-1}x_0x_1x_2x_3x_4$. Thus, for at least one of them $\{x_1, x_2, x_3, x_4\} \cap \{x_{-3}, x_{-2}\} = \emptyset$, and we may extend $Q$ by adding the vertices $x_1, x_2, x_3, x_4$ — a contradiction with the maximality of $Q$ (see Figure 6).

\begin{figure}[h]
\centering
\includegraphics[width=0.5\textwidth]{hyperpath.png}
\caption{A hyperpath $Q$ originating at $e$.}
\end{figure}

Similarly, one can prove that for most pairs of edges of $P$ there is a path of length at least $(1 - \delta^{1/4})n$ between them. This latter result has been used recently in [4] to determine asymptotically the Ramsey numbers for hypercycles. We devote to this application a separate subsection.

7.2. **Ramsey numbers for tight hypercycles.** Given a 3-uniform hypergraph $\mathcal{H}$, the *Ramsey number* $r(\mathcal{H})$ is defined as the smallest integer $N$ such that every red-blue coloring of the edges of the complete 3-uniform hypergraph $\mathcal{K}_N^{(3)}$ yields a monochromatic copy of $\mathcal{H}$. Given a suitably labeled set of vertices $\{v_1, \ldots, v_n\}$, the *tight cycle*, denoted further by $C_n^{(3)}$, has the edge set $\{v_1v_2v_3, v_2v_3v_4, v_3v_4v_5, \ldots, v_nv_1v_2\}$. The following result has been recently proved in [4].

**Theorem 7.2.** Let $\eta > 0$ be given. Then for all sufficiently large $n$,
\begin{align*}
4n - 1 \leq r(C_n^{(3)}) \leq (4 + \eta)n.
\end{align*}

*Sketch of proof of the upper bound:* Let $\mathcal{K}_N^{(3)} = \mathcal{H}_R \cup \mathcal{H}_B$, where $N \sim 4n$, be a red-blue coloring of the edges of the complete 3-uniform hypergraph $\mathcal{K}_N^{(3)}$.

We apply simultaneously to both, $\mathcal{H}_R$ and $\mathcal{H}_B$, the Hypergraph Regularity Lemma (Theorem 2.13) with suitably chosen parameters, and obtain a vertex partition $V = V_1 \cup \ldots \cup V_t$, $|V_i| = N/t$ (assume $t$ divides $N$), such that for almost all triples $\{i, j, k\}$ one of the induced
sub-hypergraphs, $\mathcal{H}_R[V_i \cup V_j \cup V_k]$ or $\mathcal{H}_B[V_i \cup V_j \cup V_k]$, contains a “well-structured” sub-sub-hypergraph, that is, a $(1/2, \delta, l, r(l), \epsilon(l))$-triad.

A modification of Proposition 7.1 yields that a well-structured hypergraph contains a long path, in our case of length almost $3N/t$, connecting every pair of typical edges of $P$, and avoiding a small set of forbidden vertices. We will build a monochromatic copy of $\mathcal{C}^{(3)}_{3n}$ mostly out of such paths, constructed within each of about $t/4$ vertex disjoint well-structured sub-hypergraphs. Thus, it is crucial to find about $t/4$ disjoint “well-structured” sub-hypergraphs in one color.

To this end, let $\mathcal{K}_R$ and $\mathcal{K}_B$ be two auxiliary hypergraphs on the vertex set $\{1, 2, \ldots, t\}$, whose edges are those triples $(i, j, k)$ for which, respectively, $\mathcal{H}_R[V_i \cup V_j \cup V_k]$ or $\mathcal{H}_B[V_i \cup V_j \cup V_k]$ contains a “well structured” sub-hypergraph. Set $\mathcal{K} = \mathcal{K}_R \cup \mathcal{K}_B$ and note that $|\mathcal{K}| \sim (1/2)$.

A substantial number of pages in [4] is devoted to proving that either $\mathcal{K}_R$ or $\mathcal{K}_B$ (say, $\mathcal{K}_R$) contains a connected matching $M$ of size $s \sim t/4$. Here “connected” means that between every two edges $e, f \in M$ there is a pseudo-path, that is, for some $p$, a sequence of not necessarily distinct edges $(e = e_1, \ldots, e_p = f)$ such that $|e_i \cap e_{i+1}| = 2$ for each $i = 1, \ldots, p - 1$.

Next, we find a long path in each sub-hypergraph $H_R[V_i, V_j, V_k]$, where $(i, j, k) \in M$. These paths are disjoint and have total length $3n - O(1)$. To connect the long paths together into a red cycle of length $3n$, we construct in $\mathcal{H}_R$ short paths (length $O(1)$) between the endpairs of long paths, being guided by the pseudo-paths linking in $\mathcal{K}_R$ the edges of $M$. In the actual proof we build the short paths first, and this is why the long paths have specified endpairs and must avoid a certain small set of vertices (to keep all paths, short or long, mutually disjoint, except for the endpairs where they meet).

### 7.3. Approximate decomposition into small diameter sub-hypergraphs.

It is easy to see that for every $n$-vertex graph and for every $\epsilon > 0$ one can partition $E(G) = E_0 \cup \cdots \cup E_k$, where $k \leq 1/\epsilon$, so that $|E_0| \leq \epsilon n^2$ and for each $1 \leq i \leq k$ the diameter of the subgraph $G_i = G[E_i]$ is at most $3/\epsilon$ (see [8]). Thus, in a sense, every dense graph can be decomposed into a bounded number of “small worlds” provided a small fraction of edges can be ignored. Here both bounds, on the diameter and on the number of subgraphs $G_i$, depend linearly on $1/\epsilon$. Using the Szemerédi Regularity Lemma [10] and Corollary 2.5(b), one may put the cap of four on the diameter, at the cost of letting $k$, the number of subgraphs in the partition, to be an enormous constant.

**Proposition 7.3.** [8] For all $\epsilon > 0$ there exist integers $K$ and $N$ such that for all $n$-vertex graphs $G$, where $n \geq N$, there is a partition $E(G) = E_0 \cup E_1 \cup \cdots \cup E_k$, where $k \leq K$, and $|E_0| \leq \epsilon n^2$, and for each $1 \leq i \leq k$, the diameter of the subgraph $G_i = G[E_i]$ is at most four.

An analogous result for 3-uniform hypergraphs follows from our Theorem 2.16 and the Hypergraph Regularity Lemma.

**Theorem 7.4.** For all $\xi > 0$ there exist integers $K$ and $N$ such that for all $n$-vertex 3-uniform hypergraphs $\mathcal{H}$, where $n \geq N$, there is a partition $\mathcal{H} = \mathcal{H}_0 \cup \cdots \cup \mathcal{H}_k$, where $k \leq K$ and $|\mathcal{H}_0| \leq \xi n^3$ and for each $1 \leq i \leq k$, every two pairs of vertices with positive degree in $\mathcal{H}_i$ are connected by a hyperpath in $\mathcal{H}_i$ of length at most twelve.
Sketch of proof: Given $\xi > 0$, set $t_0 = 8/\xi$, $\alpha = \xi/8$ and let $\delta_0 > 0$ and functions $r(l)$, $\epsilon(l)$ be as in Theorem 2.16 with above $\alpha$. Further, let $N_1$ be the smallest natural number, for which Theorem 2.16 holds. Set

$$\delta = \min \left\{ \delta_0, \left( \frac{\xi}{16} \right)^2 \right\}$$

and apply the Hypergraph Regularity Lemma (Theorem 2.13) with the above $\delta$, $\epsilon(l)$ and $r(l, t) = r(l)$, to get $T_0, L_0$, and $N_0$. Set

$$K = \left( \frac{T_0}{3} \right) L_0^3 \quad \text{and} \quad N = \max \left\{ N_0, \frac{16T_0}{\xi}, N_1 T_0 \right\}$$

and let $\mathcal{H}$ be an arbitrary 3-uniform hypergraph with $n \geq N$ vertices and $|\mathcal{H}| \geq \xi n^3$ triplets.

Let $(\mathcal{H}', P_s)$, $s = 1, \ldots, k \leq \binom{t}{3} < K$, be all $(\geq \alpha, \delta, l, r(l), \epsilon(l))$-triads $(\mathcal{H}', P)$, where $\mathcal{H}' = \mathcal{H} \cap \text{Tr}(P)$ and $P = (P_{ai}^h, P_{bj}^h, P_{ci}^h)$, $1 \leq i < j < h \leq t$, $1 \leq a, b, c \leq l$, resulting from applying Theorem 2.13 to $\mathcal{H}$. For each $s = 1, \ldots, k$, let $(P_s)_0$ be the subgraph of $P_s$ guaranteed by Theorem 2.16, and set $\mathcal{H}_s = \mathcal{H}'_s - (P_s)_0$. Then, each pair of edges of $P_s - (P_s)_0$, that is each pair of edges of $P_s$ with positive degree in $\mathcal{H}_s$, is connected in $\mathcal{H}_s$ by a hyperpath of length at most twelve.

Let us set $\mathcal{H}_0 = \mathcal{H} \setminus \bigcup_{s=1}^k \mathcal{H}_s$. To complete the proof of Theorem 7.4, it remains to show that $|\mathcal{H}_0| \leq \xi n^3$. The edges that belong to $\mathcal{H}_0$ either intersect the set $V_0$, or intersect a set $V_i$, $i \geq 1$, in more than one vertex, or belong to $(\delta, r)$-irregular triads, or to triads with density less than $\alpha$. We omit the details of tedious but straightforward calculations. \hfill \Box

References


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